

How to
manipulate curve standards:
a white paper for the black hat

Daniel J. Bernstein

Tung Chou

Chitchanok Chuengsatiansup

Andreas Hülsing

Eran Lambooj

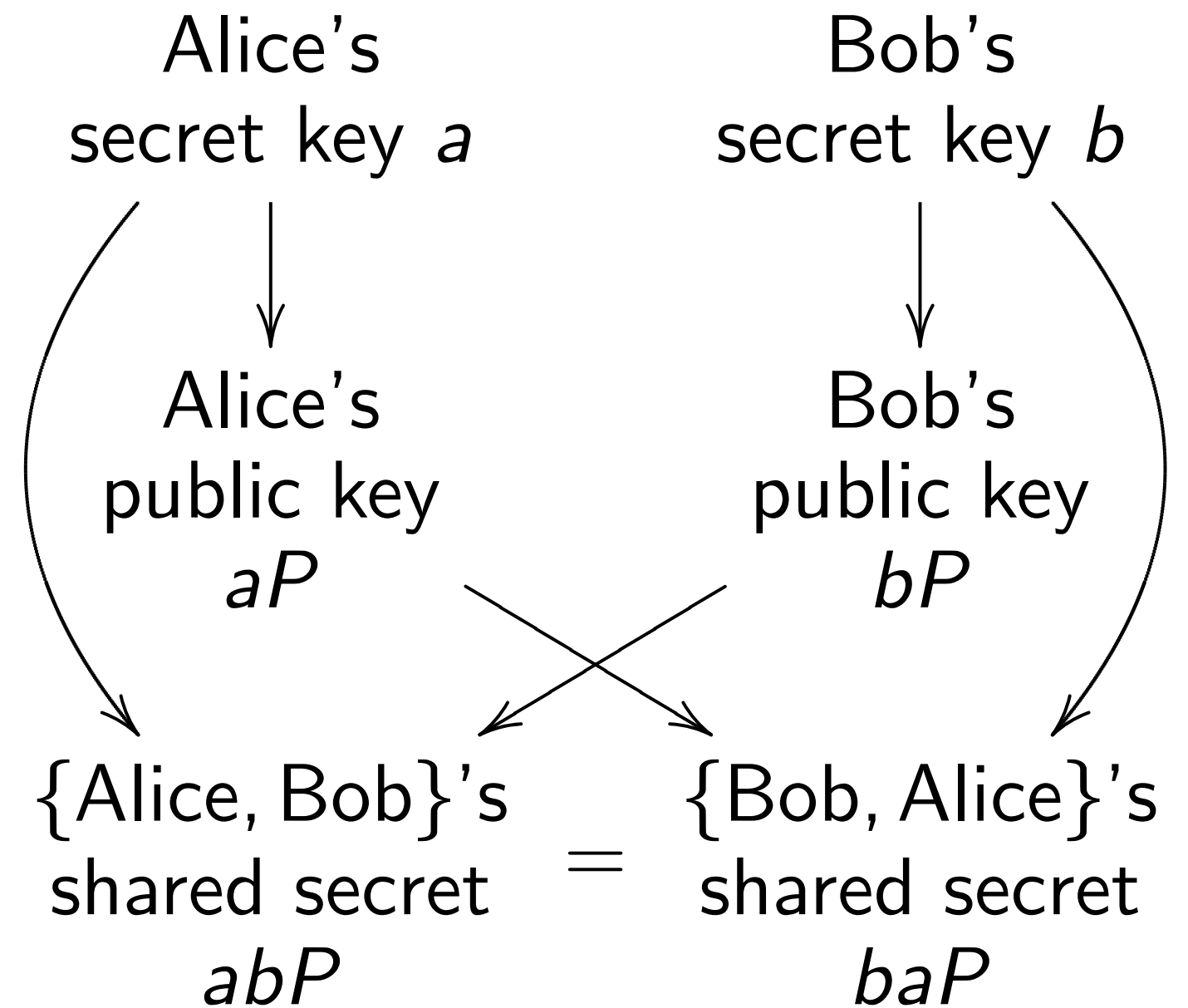
Tanja Lange

Ruben Niederhagen

Christine van Vredendaal

bada55.cr.yp.to

Textbook key exchange
using standard point P
on a standard elliptic curve E :



How to
manipulate curve standards:
a white paper for the black hat

Daniel J. Bernstein

Tung Chou

Chitchanok Chuengsatiansup

Andreas Hülsing

Eran Lambooj

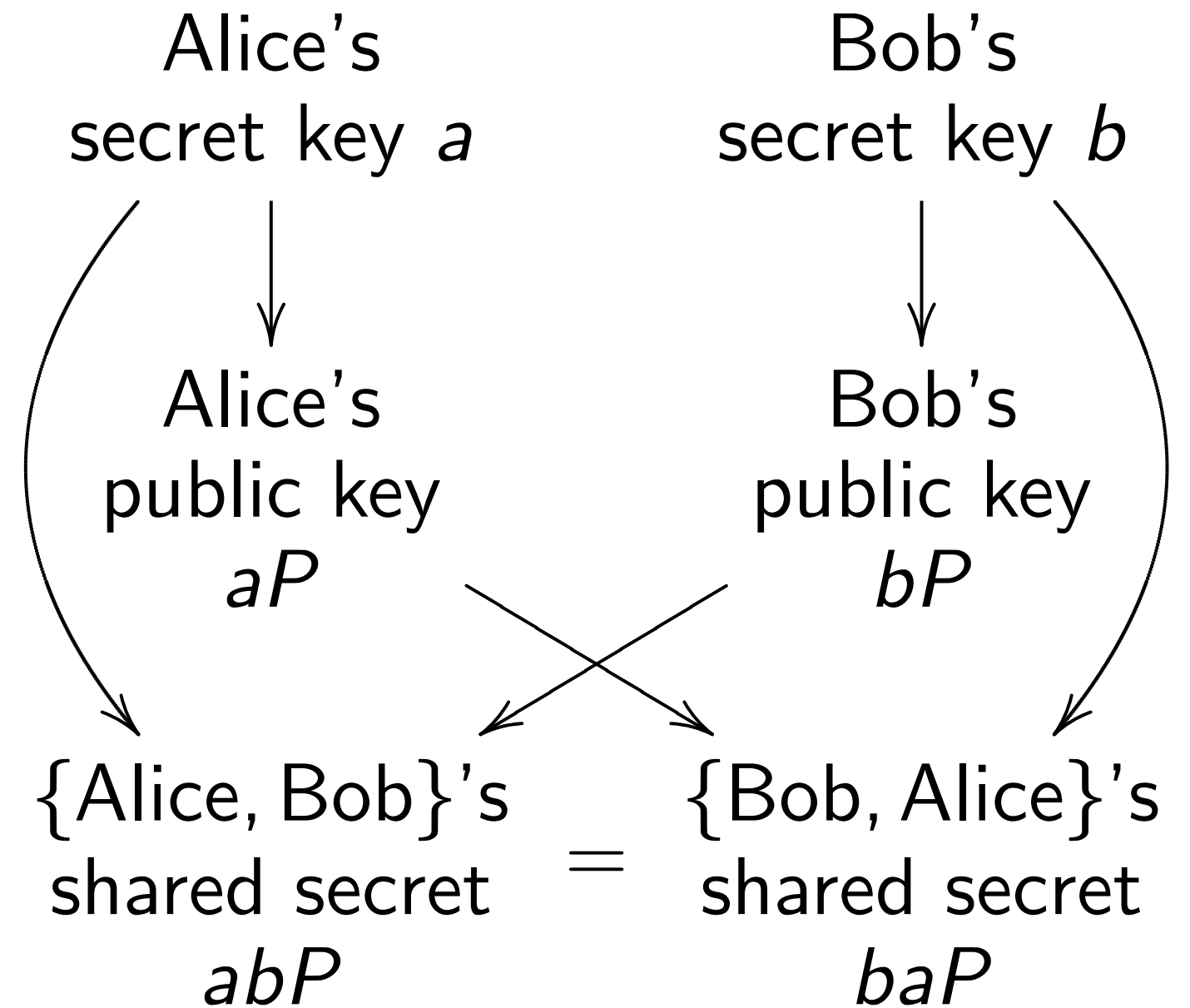
Tanja Lange

Ruben Niederhagen

Christine van Vredendaal

bada55.cr.yp.to

Textbook key exchange
using standard point P
on a standard elliptic curve E :



Security depends on choice of E .

ate curve standards:
paper for the black hat

. Bernstein

hou

ok Chuengsatiansup

Hülsing

mbooij

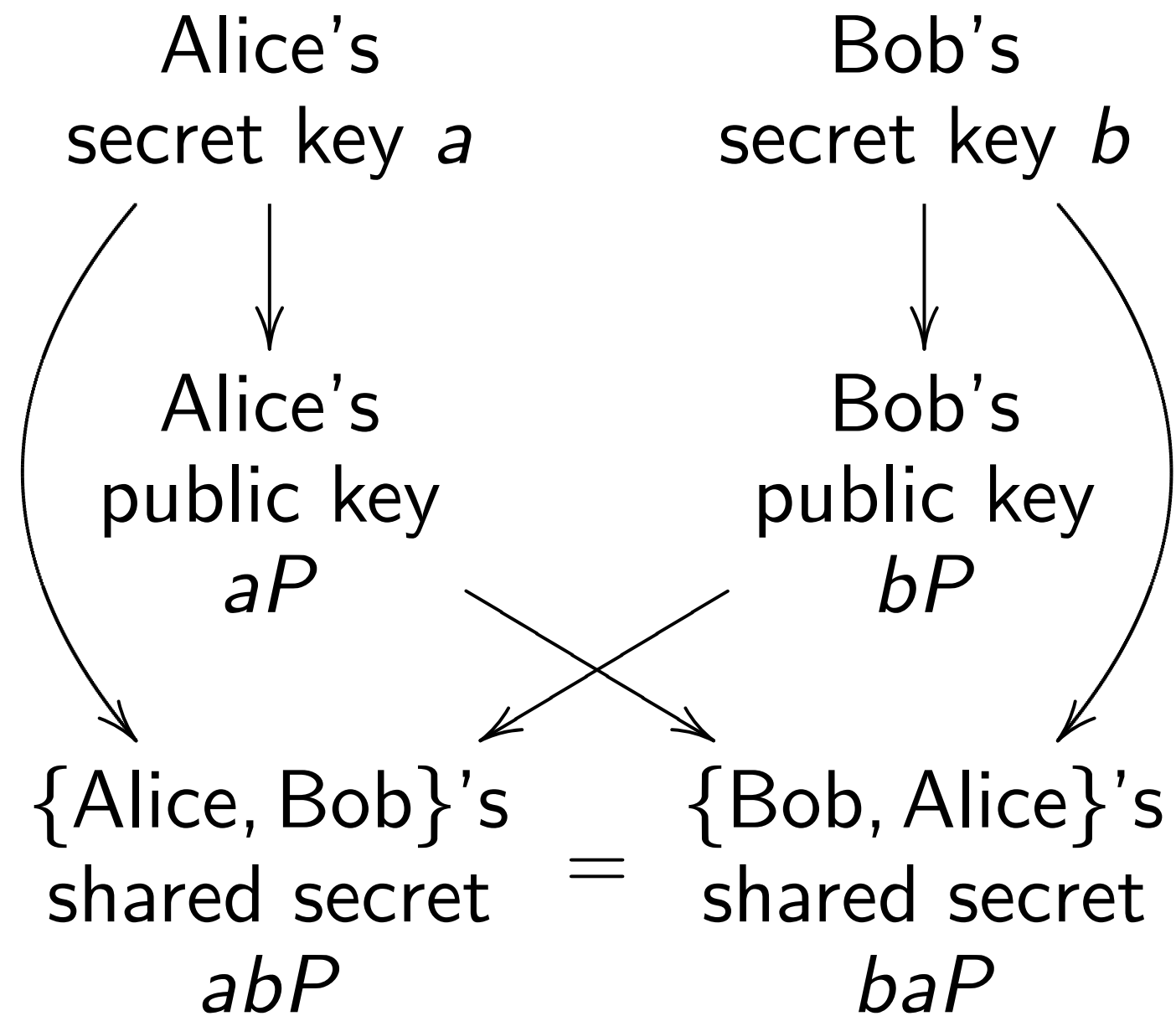
ange

Niederhagen

e van Vredendaal

cr.yp.to

Textbook key exchange
using standard point P
on a standard elliptic curve E :



Security depends on choice of E .

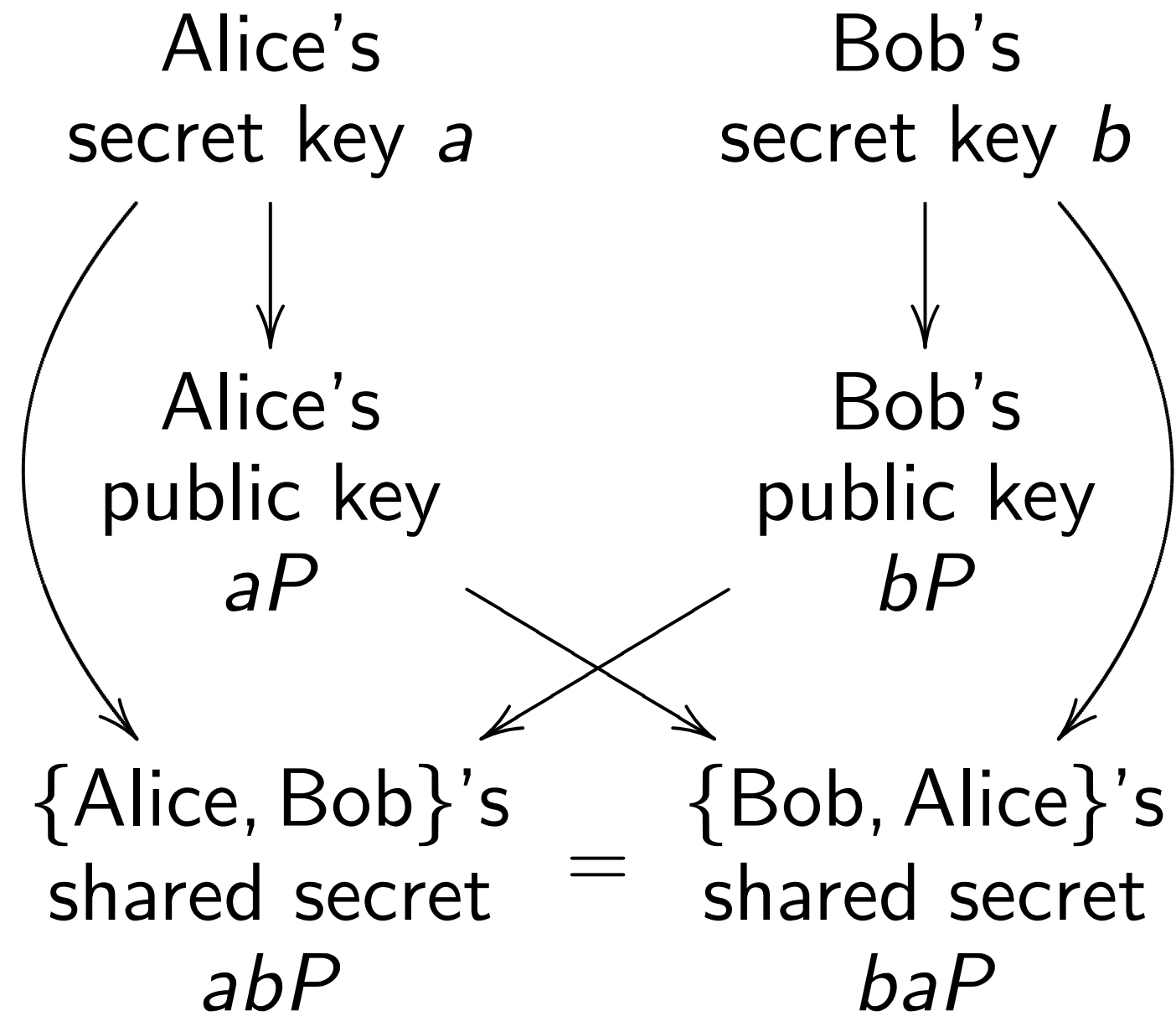


standards:
the black hat

n
gsatiansup

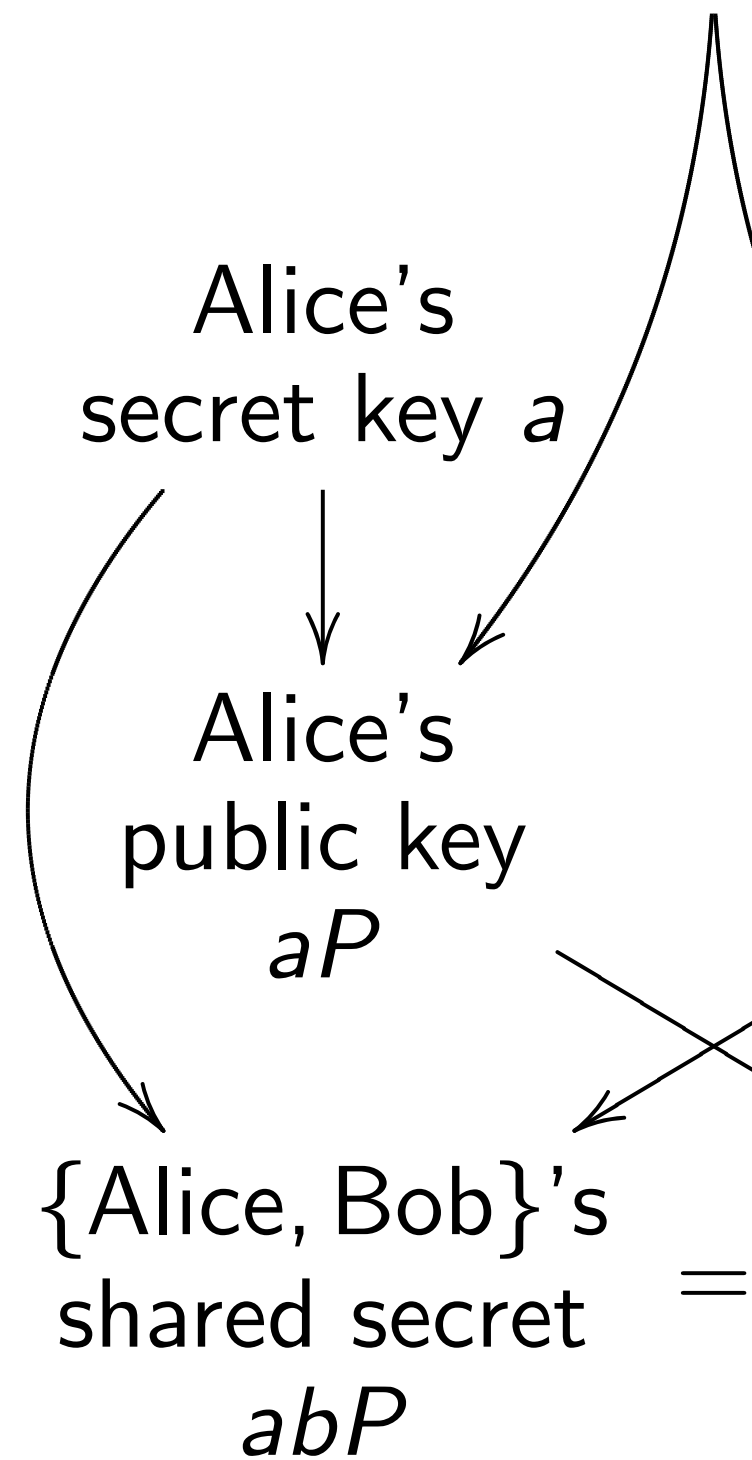
n
lendaal

Textbook key exchange
using standard point P
on a standard elliptic curve E :

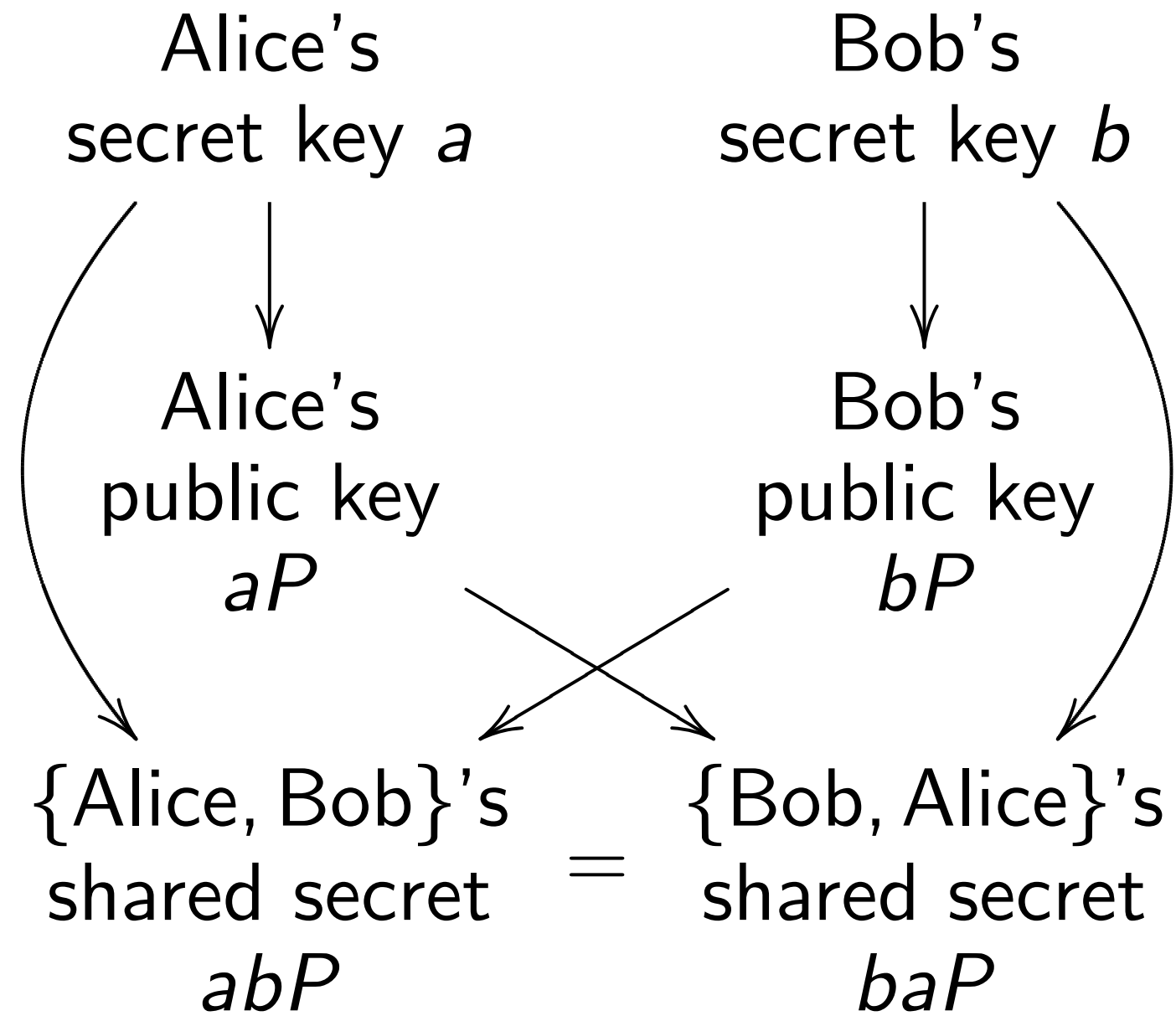


Security depends on choice of E .

Our partner
choice o

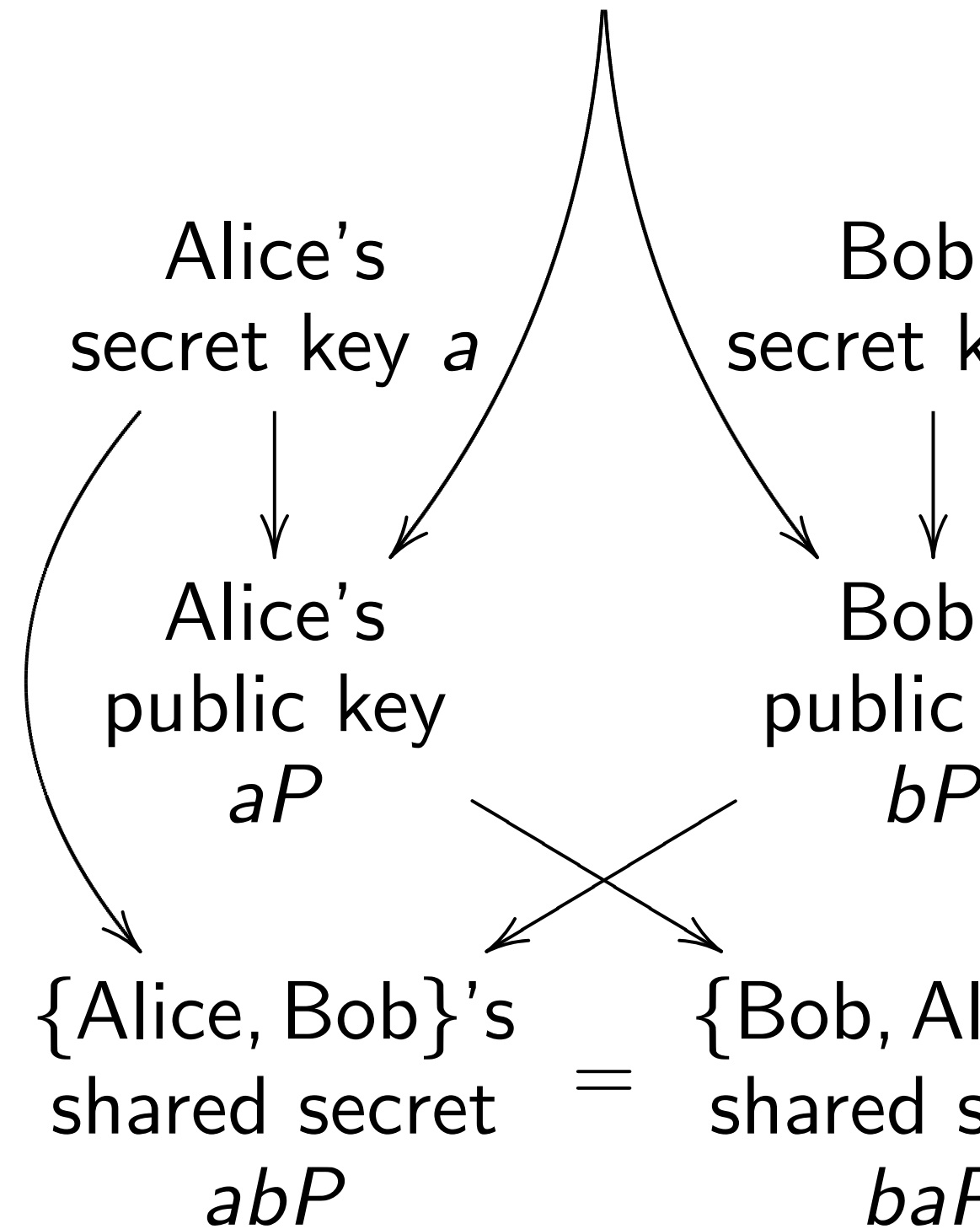


Textbook key exchange
using standard point P
on a standard elliptic curve E :

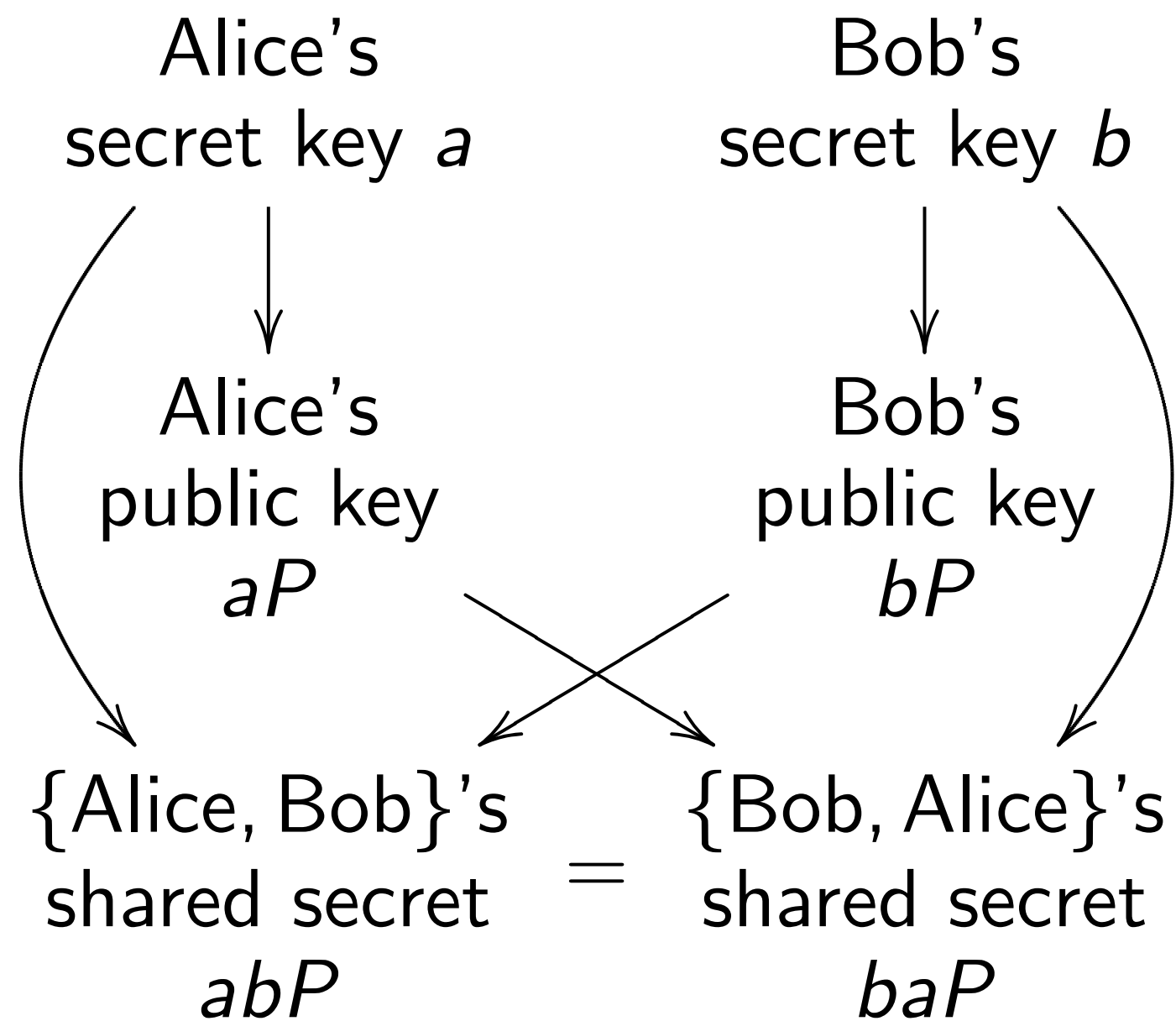


Security depends on choice of E .

Our partner Jerry's
choice of E, P

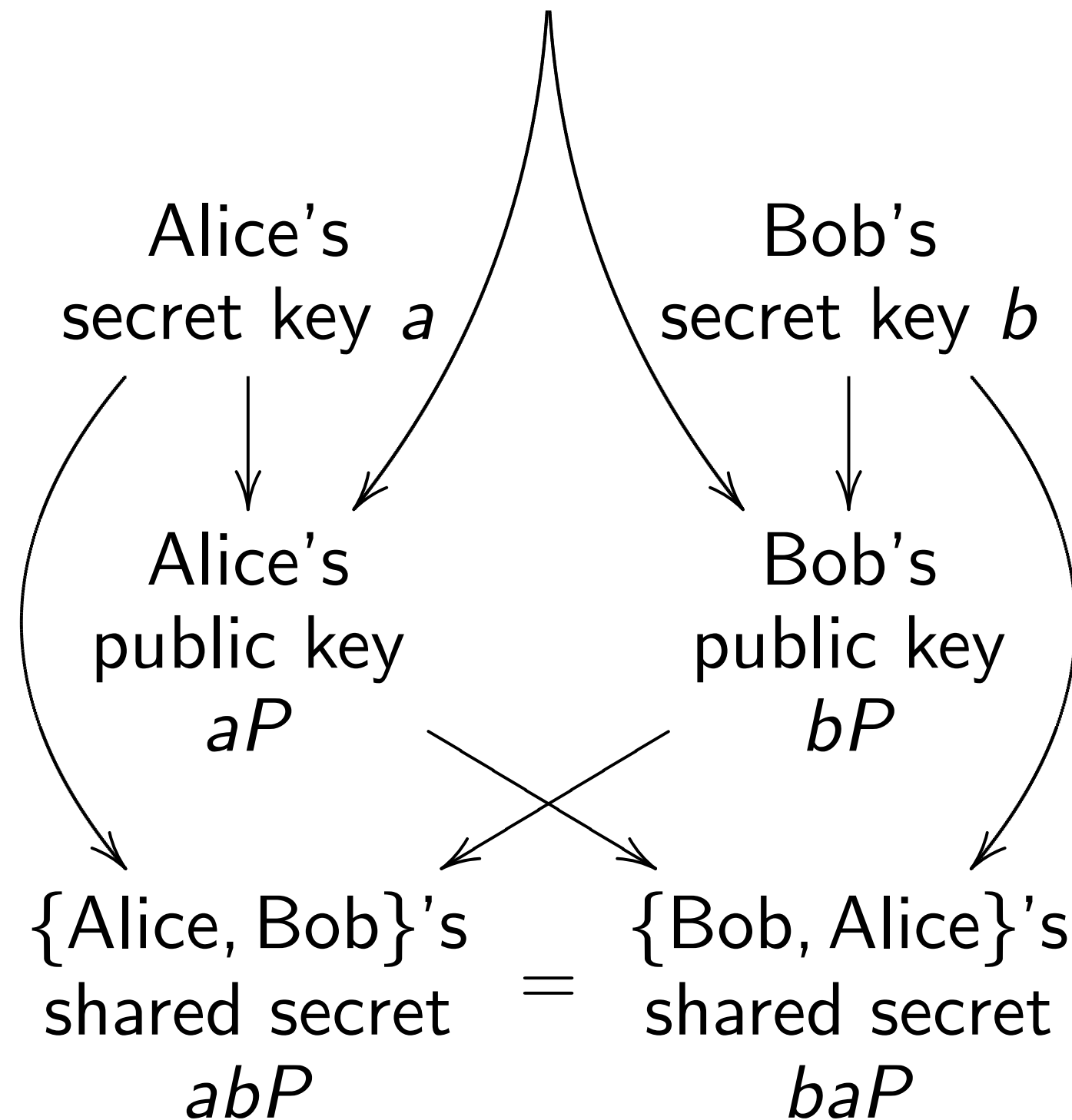


Textbook key exchange
 using standard point P
 on a standard elliptic curve E :

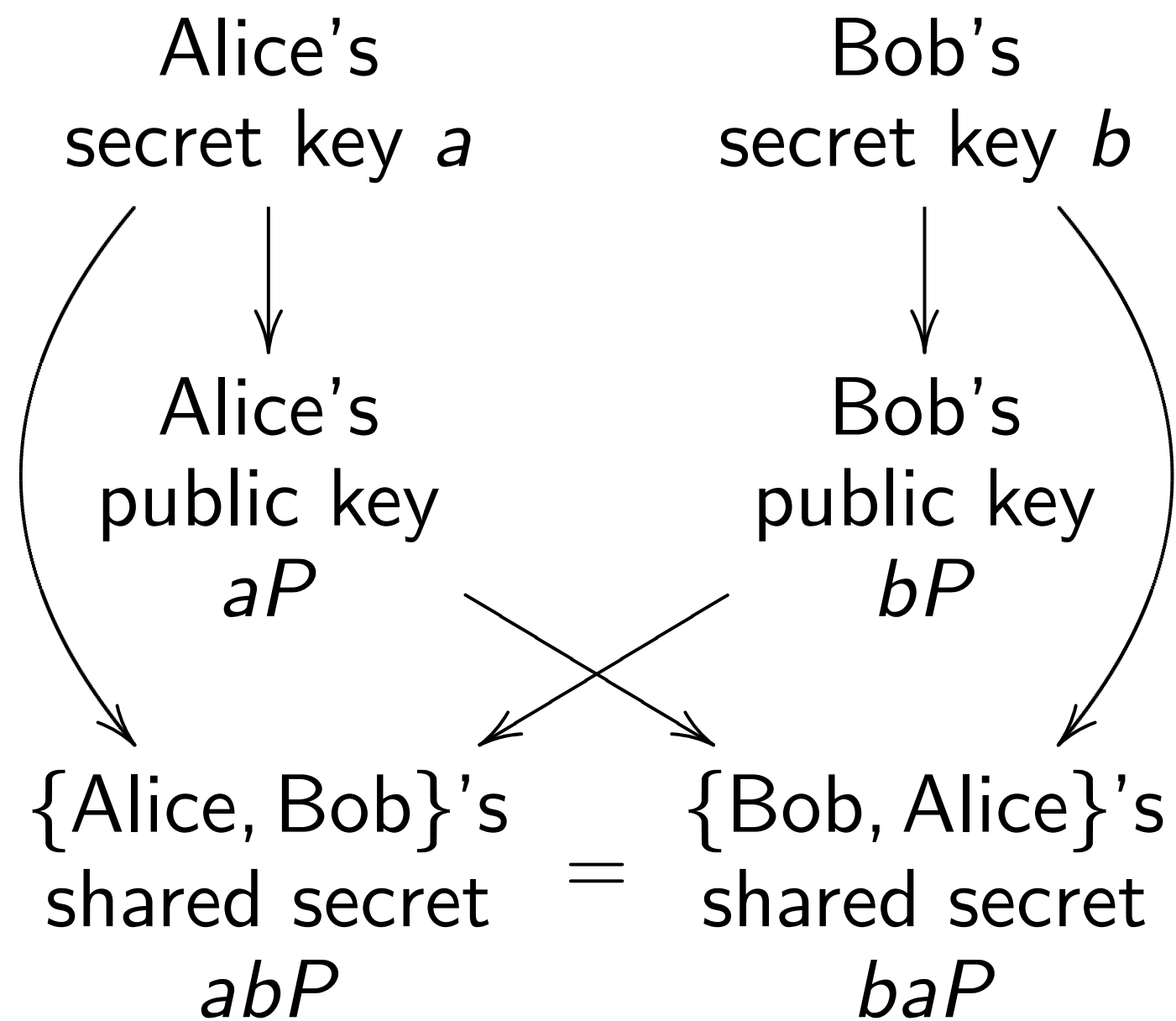


Security depends on choice of E .

Our partner Jerry's
 choice of E, P

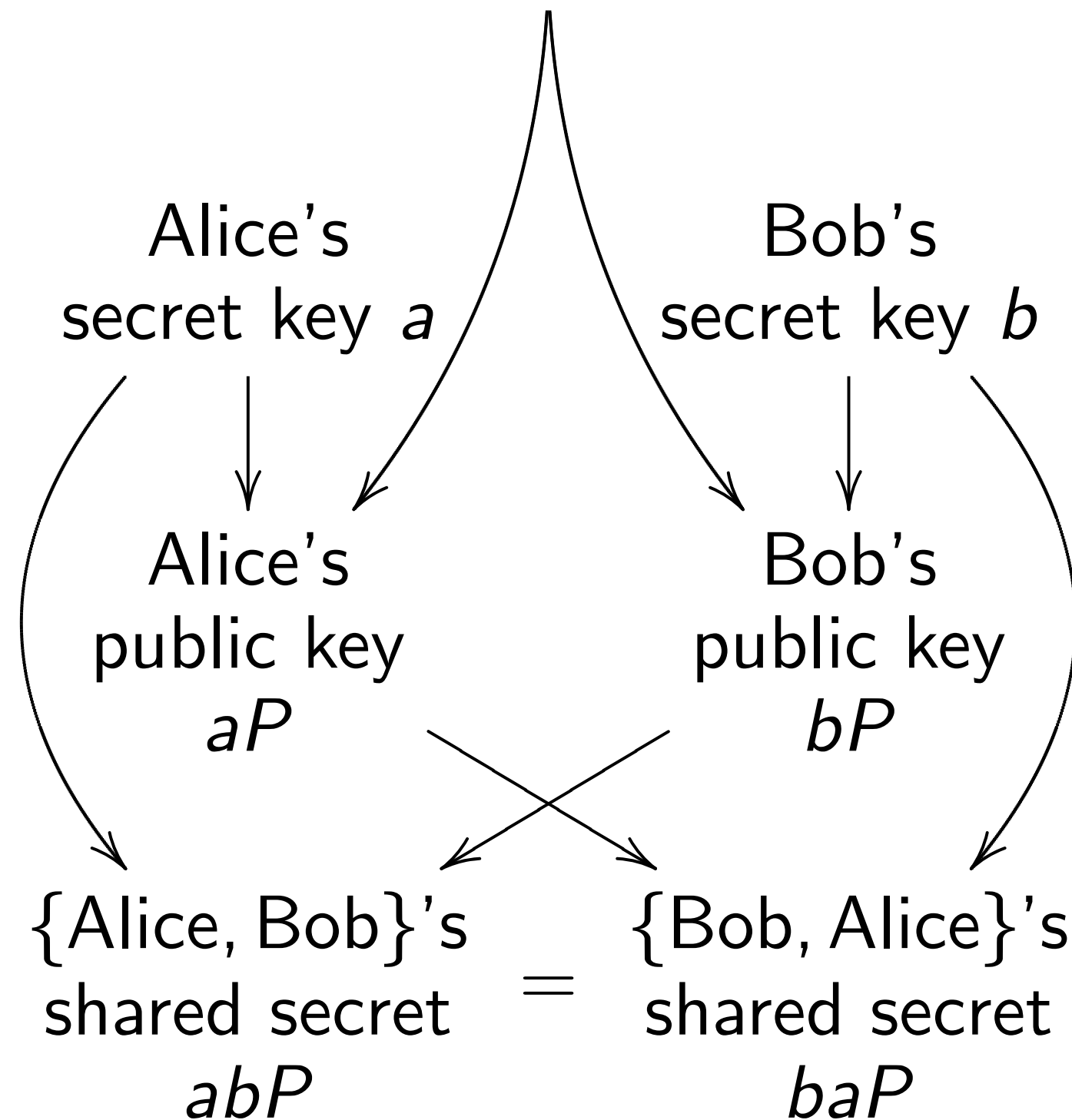


Textbook key exchange
using standard point P
on a standard elliptic curve E :



Security depends on choice of E .

Our partner Jerry's
choice of E, P

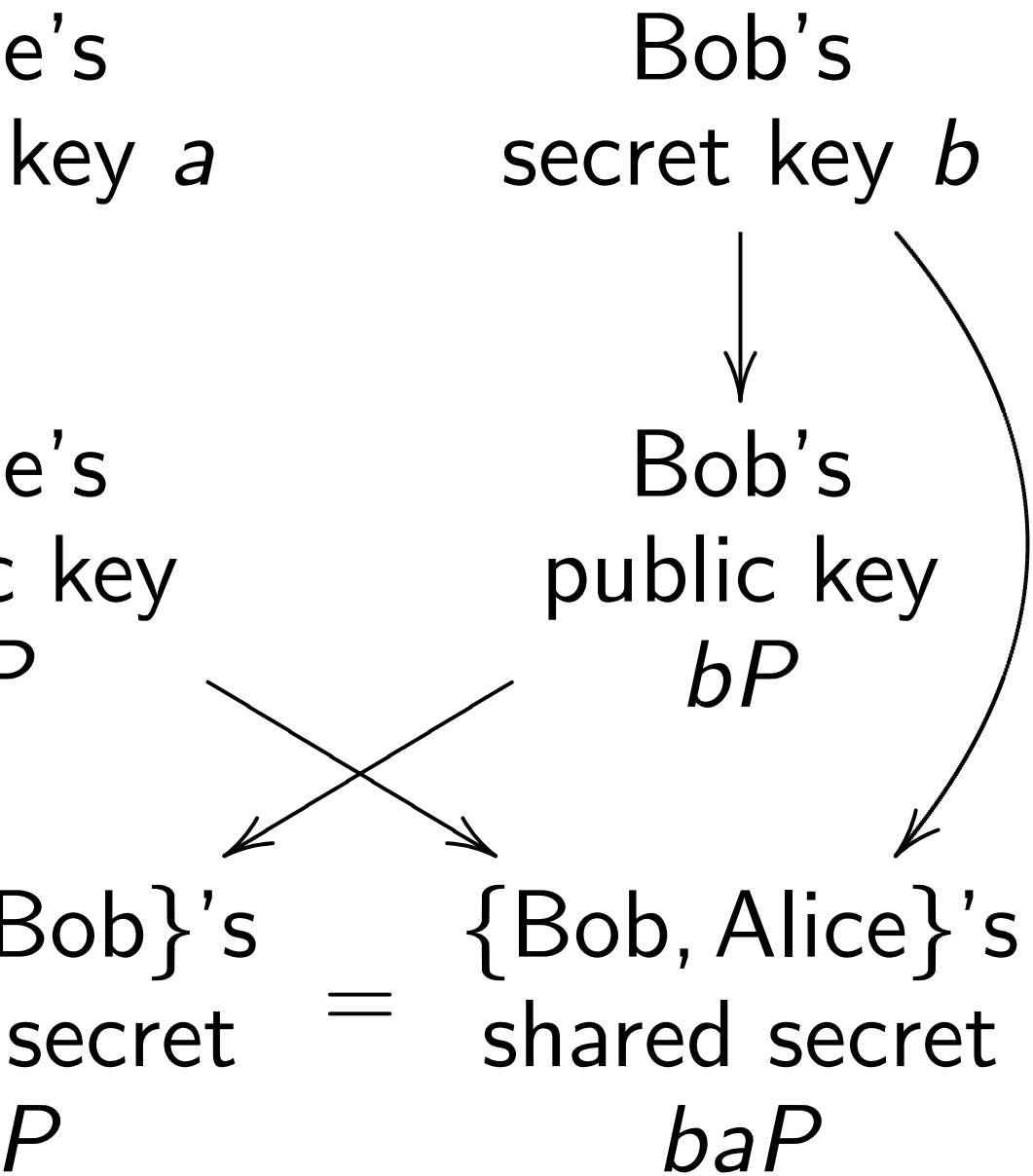


Can we exploit this picture?

key exchange

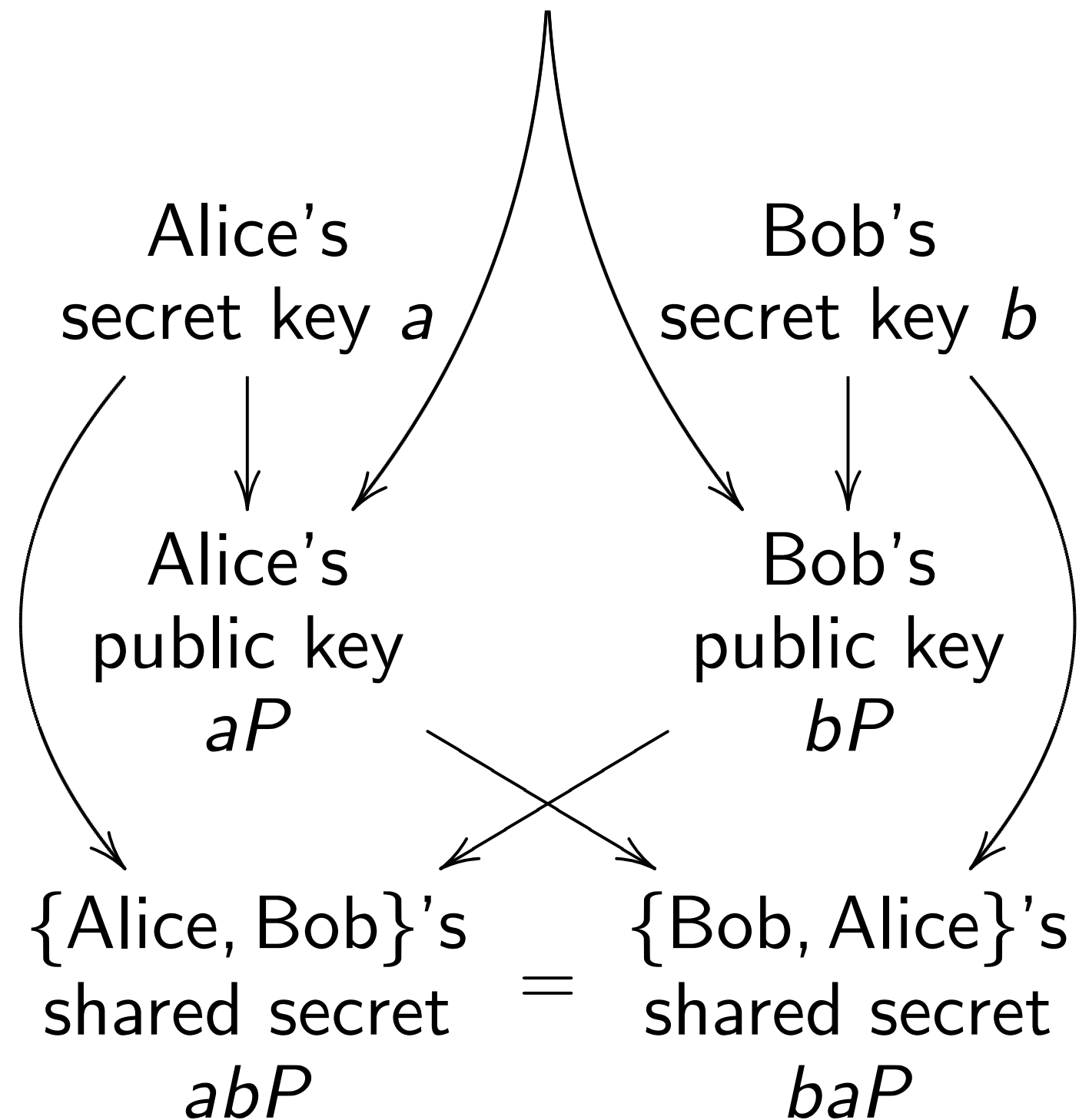
standard point P

standard elliptic curve E :



depends on choice of E .

Our partner Jerry's choice of E, P



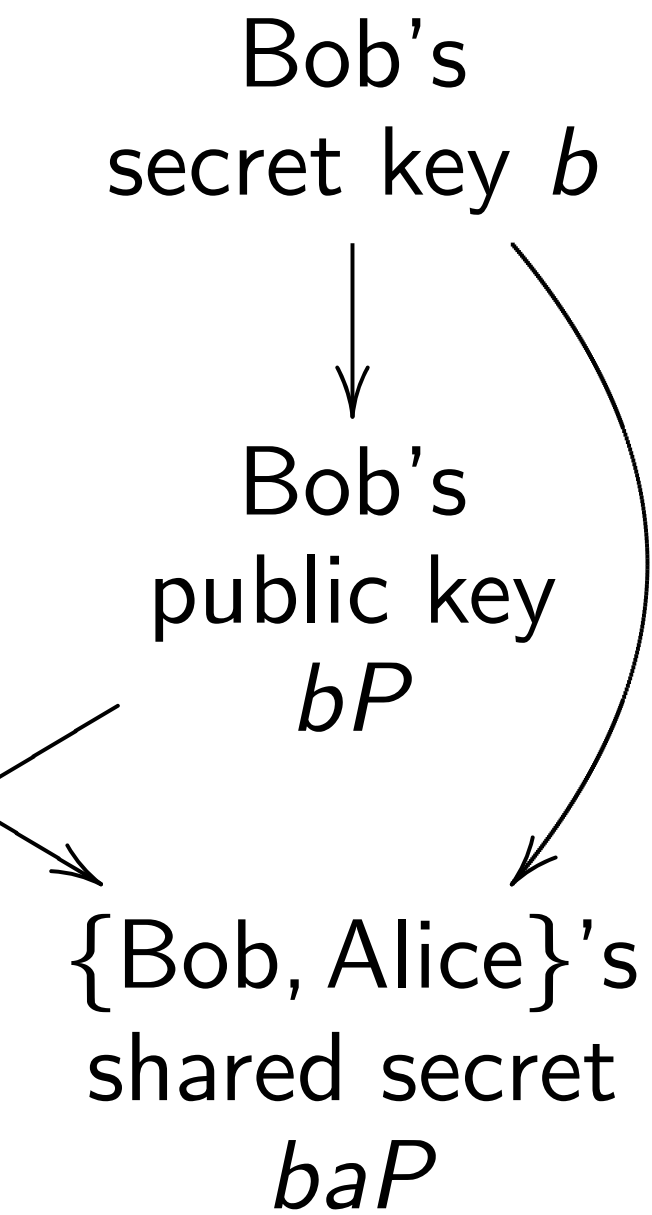
Can we exploit this picture?

Exploita
public cr

change

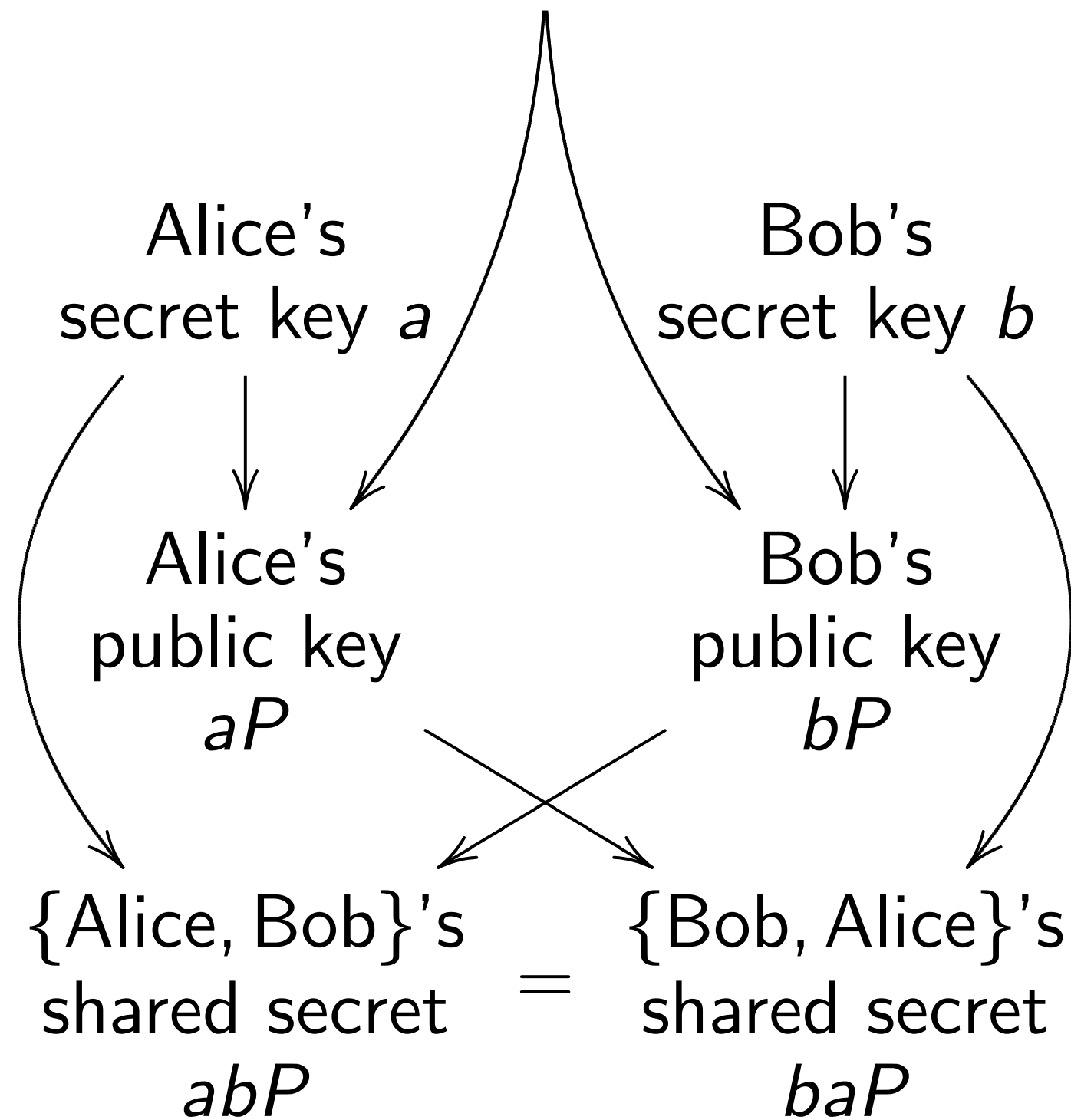
point P

elliptic curve E :



on choice of E .

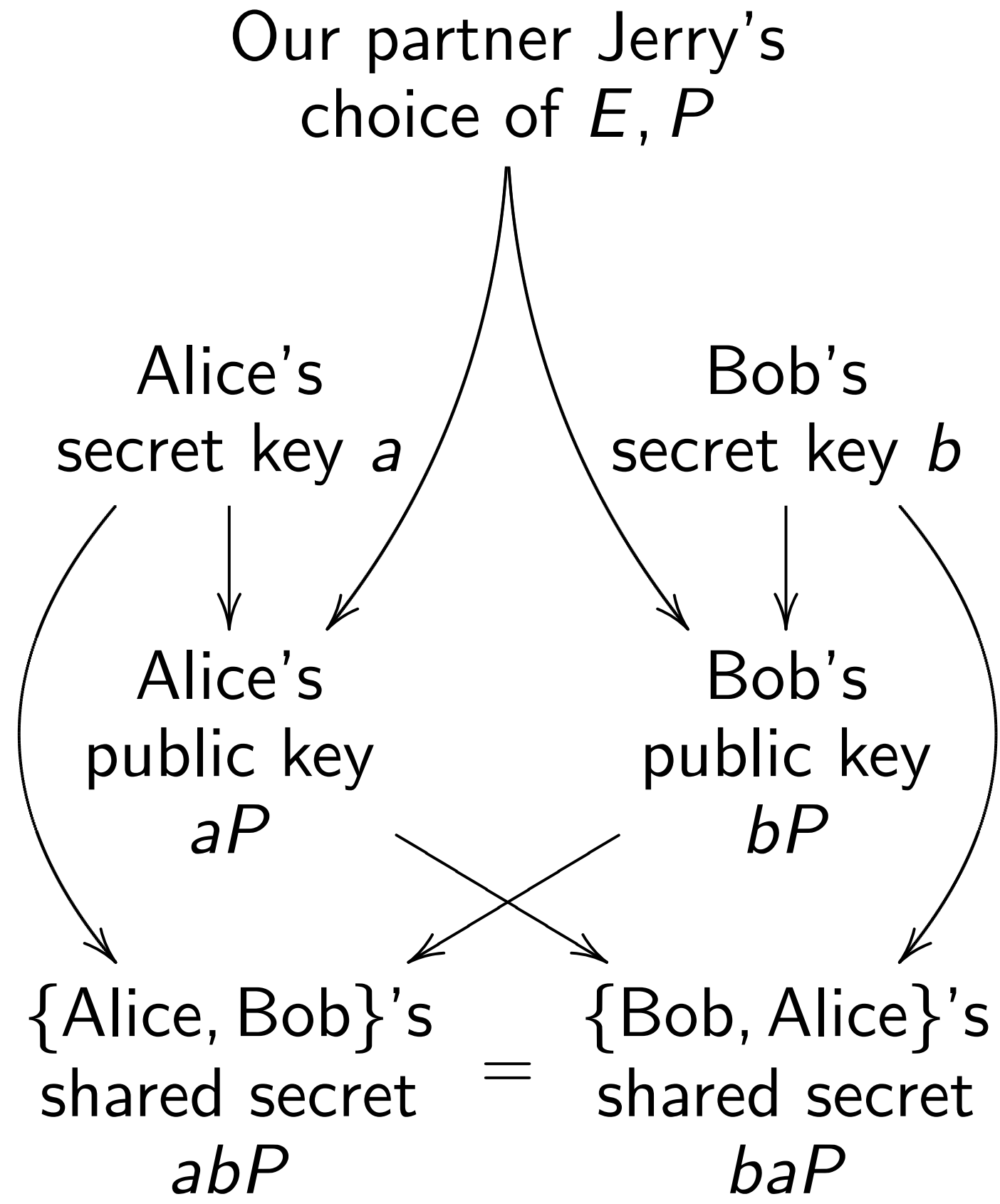
Our partner Jerry's choice of E, P



Can we exploit this picture?

Exploitability depends on public criteria for a

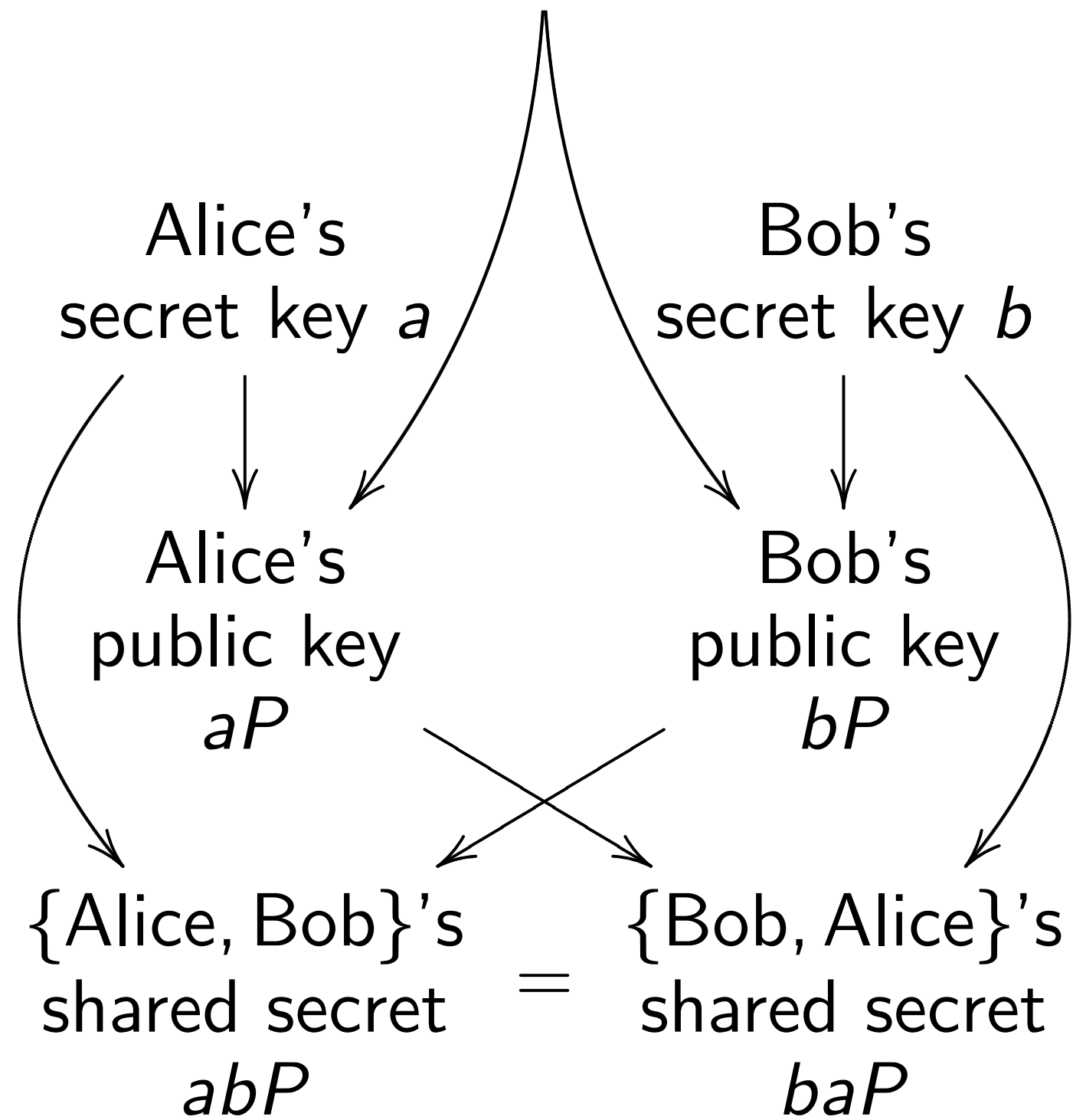
E :
Alice's secret key b
Alice's public key
Alice's secret key
of E .



Can we exploit this picture?

Exploitability depends on public criteria for accepting

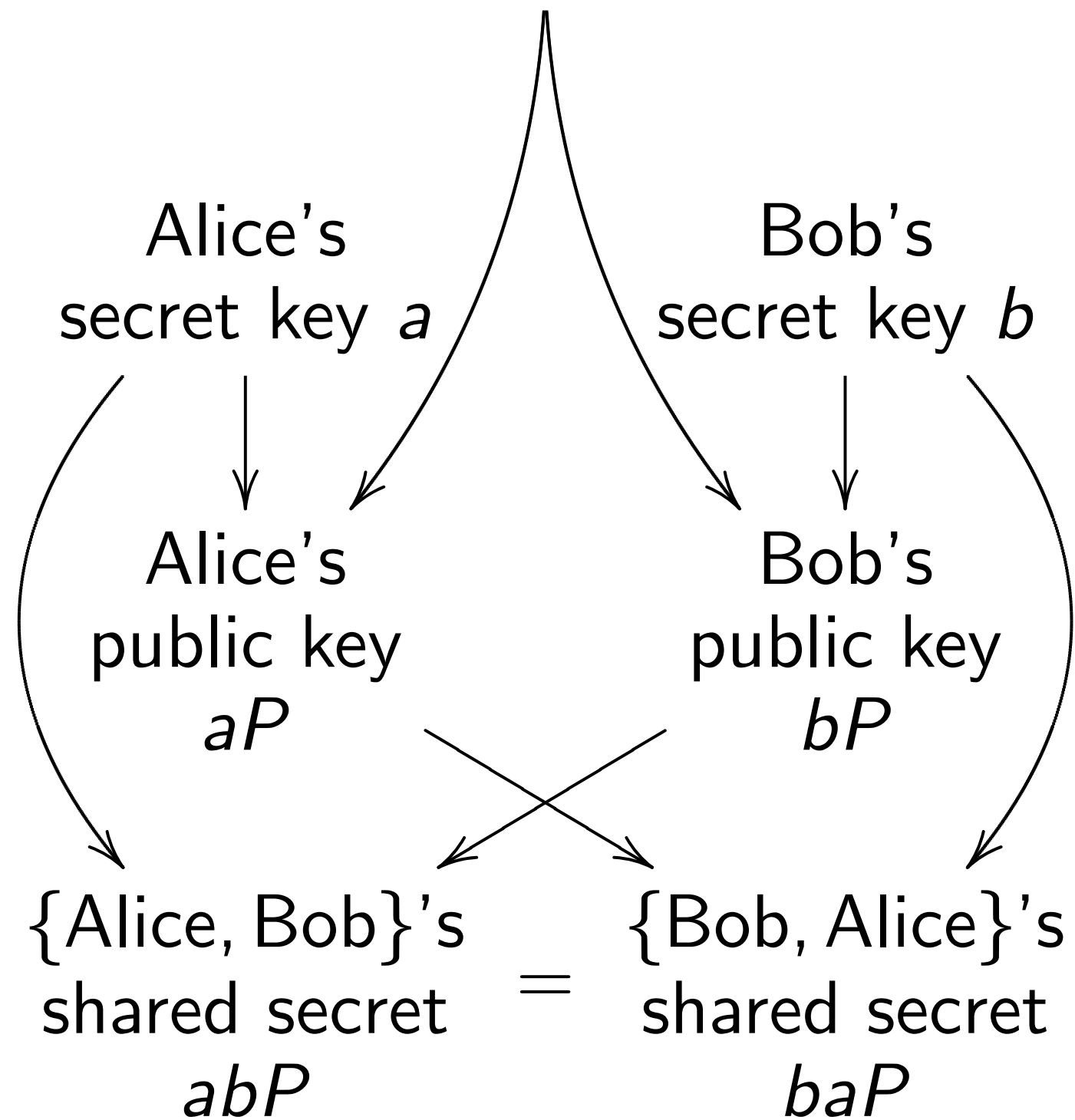
Our partner Jerry's
choice of E, P



Exploitability depends on
public criteria for accepting E, P .

Can we exploit this picture?

Our partner Jerry's
choice of E, P



Can we exploit this picture?

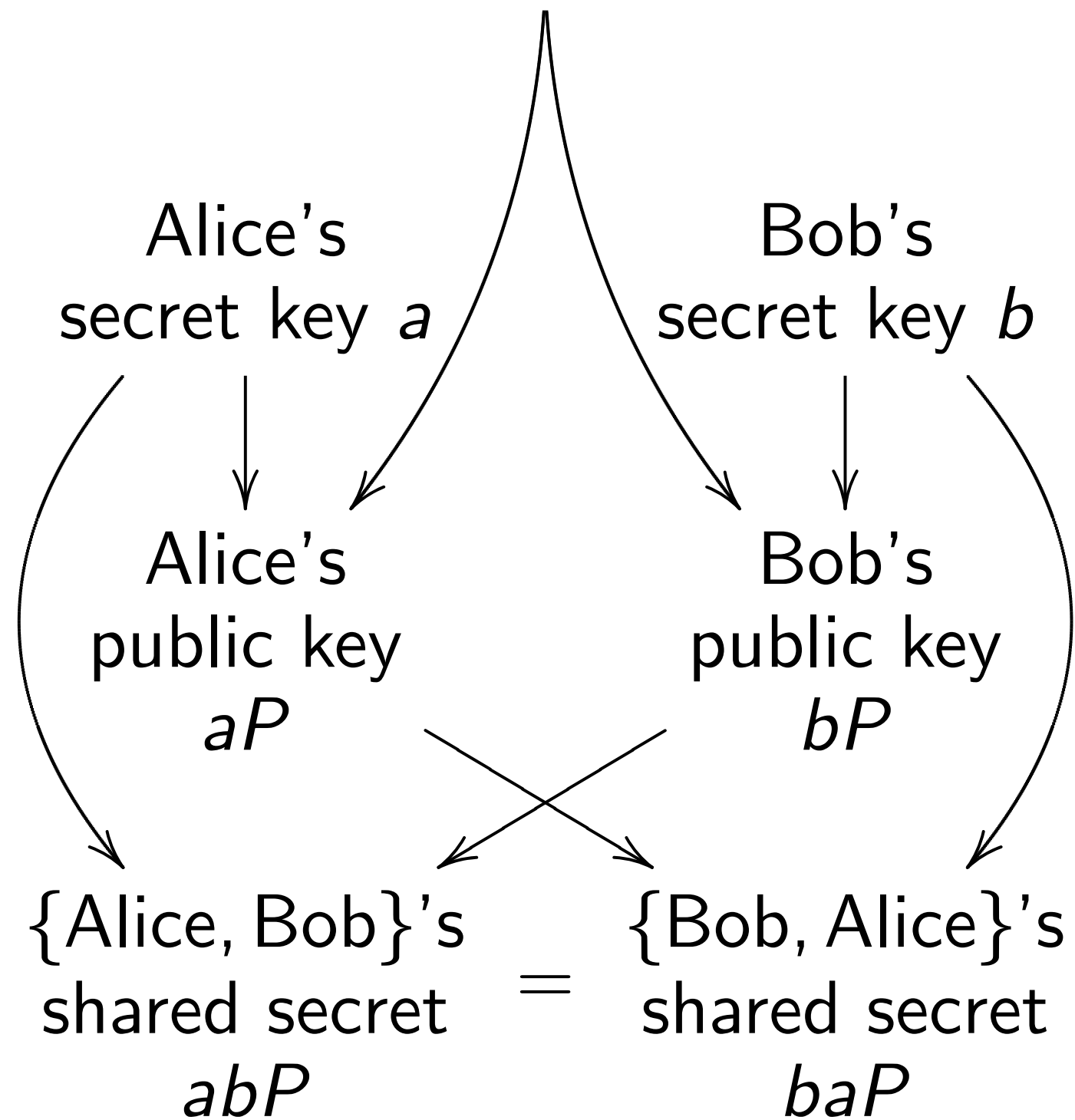
Exploitability depends on
public criteria for accepting E, P .

Extensive ECC literature:

Pollard rho breaks small E ,
Pohlig–Hellman breaks most E ,
MOV/FR breaks some E ,
SmartASS breaks some E , etc.

Assume that public will accept
any E not publicly broken.

Our partner Jerry's
choice of E, P



Can we exploit this picture?

Exploitability depends on
public criteria for accepting E, P .

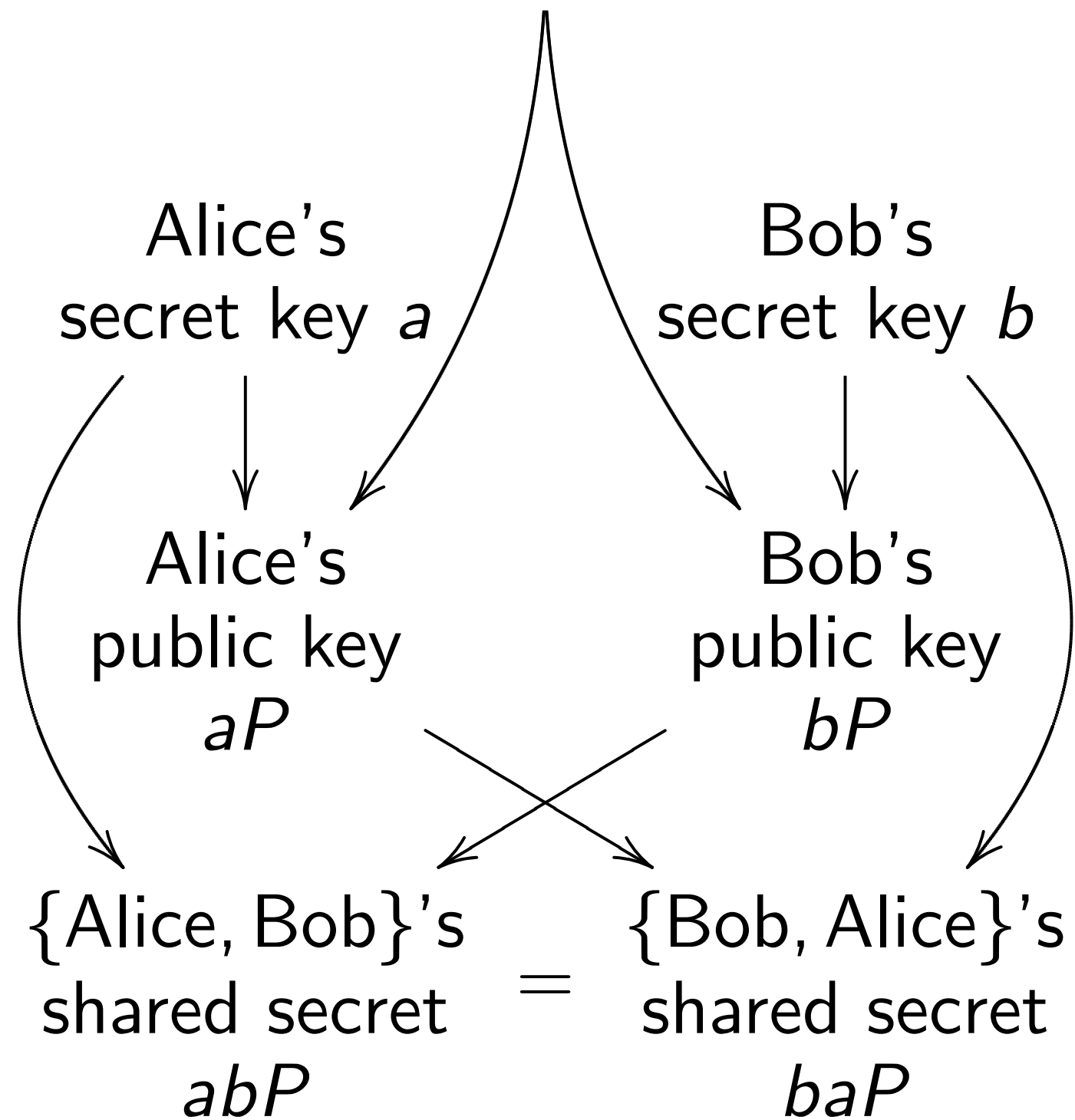
Extensive ECC literature:

Pollard rho breaks small E ,
Pohlig–Hellman breaks most E ,
MOV/FR breaks some E ,
SmartASS breaks some E , etc.

Assume that public will accept
any E not publicly broken.

Assume that we've figured out
how to break another curve E .

Our partner Jerry's
choice of E, P



Can we exploit this picture?

Exploitability depends on
public criteria for accepting E, P .

Extensive ECC literature:

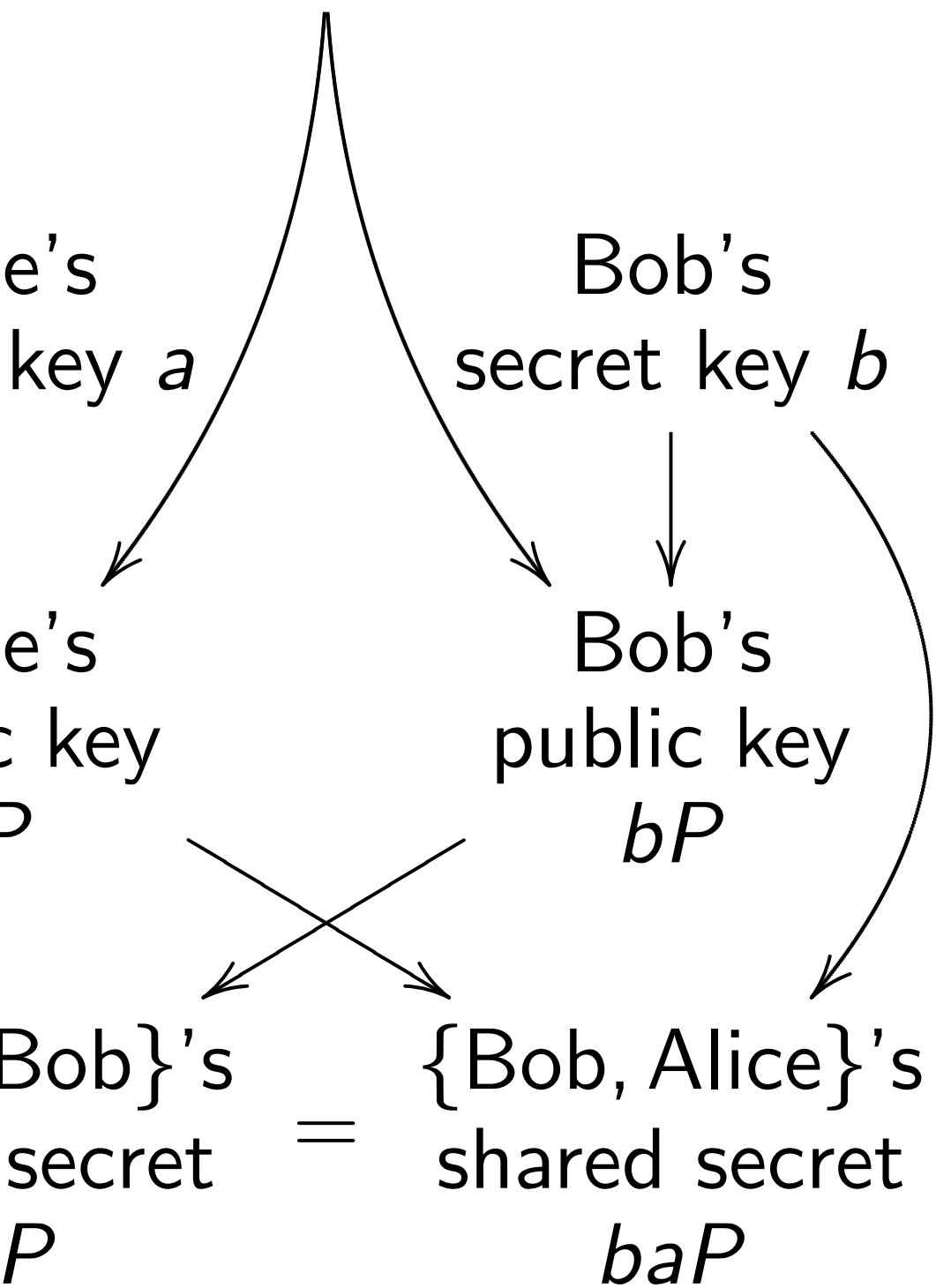
Pollard rho breaks small E ,
Pohlig–Hellman breaks most E ,
MOV/FR breaks some E ,
SmartASS breaks some E , etc.

Assume that public will accept
any E not publicly broken.

Assume that we've figured out
how to break another curve E .

Jerry standardizes this curve.
Alice and Bob use it.

Our partner Jerry's
choice of E, P



exploit this picture?

Exploitability depends on
public criteria for accepting E, P .

[Extensive ECC literature:](#)

Pollard rho breaks small E ,
Pohlig–Hellman breaks most E ,
MOV/FR breaks some E ,
SmartASS breaks some E , etc.

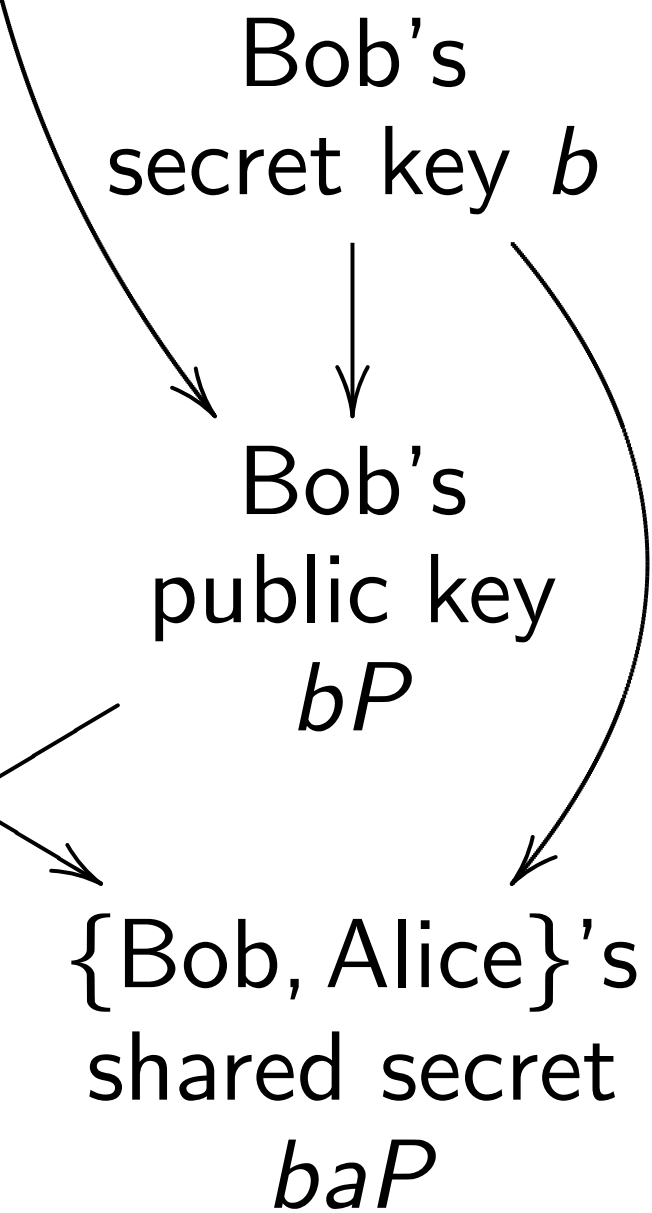
Assume that public will accept
any E not publicly broken.

Assume that we've figured out
how to break another curve E .

Jerry standardizes this curve.
Alice and Bob use it.

Is first a
Would t
any curv
that surv

er Jerry's
f E, P



s picture?

Exploitability depends on
public criteria for accepting E, P .

Extensive ECC literature:

Pollard rho breaks small E ,
Pohlig–Hellman breaks most E ,
MOV/FR breaks some E ,
SmartASS breaks some E , etc.

Assume that public will accept
any E not publicly broken.

Assume that we've figured out
how to break another curve E .


Jerry standardizes this curve.
Alice and Bob use it.

Is first assumption
Would the public
any curve chosen
that survives these

's
key b

's
key

Alice}'s
secret
b



Exploitability depends on public criteria for accepting E, P .

Extensive ECC literature:

Pollard rho breaks small E ,
Pohlig–Hellman breaks most E ,
MOV/FR breaks some E ,
SmartASS breaks some E , etc.

Assume that public will accept any E not publicly broken.

Assume that we've figured out how to break another curve E .

Jerry standardizes this curve.

Alice and Bob use it.

Is first assumption plausible?

Would the public really accept *any* curve chosen by Jerry that survives these criteria?

Exploitability depends on public criteria for accepting E, P .

Extensive ECC literature:

Pollard rho breaks small E ,
Pohlig–Hellman breaks most E ,
MOV/FR breaks some E ,
SmartASS breaks some E , etc.

Assume that public will accept any E not publicly broken.

Assume that we've figured out how to break another curve E .

Jerry standardizes this curve.

Alice and Bob use it.

Is first assumption plausible?

Would the public really accept *any* curve chosen by Jerry that survives these criteria?

Exploitability depends on public criteria for accepting E, P .

Extensive ECC literature:

Pollard rho breaks small E ,

Pohlig–Hellman breaks most E ,

MOV/FR breaks some E ,

SmartASS breaks some E , etc.

Assume that public will accept any E not publicly broken.

Assume that we've figured out how to break another curve E .

Jerry standardizes this curve.

Alice and Bob use it.

Is first assumption plausible?

Would the public really accept *any* curve chosen by Jerry that survives these criteria?

Example showing plausibility:
Chinese OSCCA SM2 (2010) includes algorithms and a curve. The curve looks random; survives these criteria; has no other justification.

Exploitability depends on public criteria for accepting E, P .

Extensive ECC literature:

Pollard rho breaks small E ,
Pohlig–Hellman breaks most E ,
MOV/FR breaks some E ,
SmartASS breaks some E , etc.

Assume that public will accept any E not publicly broken.

Assume that we've figured out how to break another curve E .

Jerry standardizes this curve.

Alice and Bob use it.

Is first assumption plausible?

Would the public really accept *any* curve chosen by Jerry that survives these criteria?

Example showing plausibility:
Chinese OSCCA SM2 (2010) includes algorithms and a curve. The curve looks random; survives these criteria; has no other justification.

More recent example:
French [ANSSI FRP256V1](#) (2011).
Again no justification.

bility depends on
criteria for accepting E, P .

[the ECC literature](#):

who breaks small E ,
Hellman breaks most E ,
R breaks some E ,
SS breaks some E , etc.

that public will accept
not publicly broken.

that we've figured out
break another curve E .

standardizes this curve.

and Bob use it.

Is first assumption plausible?

Would the public really accept
any curve chosen by Jerry
that survives these criteria?

Example showing plausibility:
Chinese OSCCA SM2 (2010)
includes algorithms and a curve.
The curve looks random;
survives these criteria;
has no other justification.

More recent example:
French [ANSSI FRP256V1](#) (2011).
Again no justification.

Maybe p
outside
 E must
and Jerry
"seed" s

ends on
accepting E, P .

erature:

small E ,
breaks most E ,
some E ,
some E , etc.

c will accept
broken.

e figured out
her curve E .

this curve.
it.

Is first assumption plausible?

Would the public really accept
any curve chosen by Jerry
that survives these criteria?

Example showing plausibility:
Chinese OSCCA SM2 (2010)
includes algorithms and a curve.
The curve looks random;
survives these criteria;
has no other justification.

More recent example:
French [ANSSI FRP256V1](#) (2011).
Again no justification.

Maybe public is more
outside China and
 E must not be pu
and Jerry must pr
“seed” s such that

Is first assumption plausible?

Would the public really accept *any* curve chosen by Jerry that survives these criteria?

Example showing plausibility:
Chinese OSCCA SM2 (2010)
includes algorithms and a curve.
The curve looks random;
survives these criteria;
has no other justification.

More recent example:
French [ANSSI FRP256V1](#) (2011).
Again no justification.

Maybe public is more demanding
outside China and France:
E must not be publicly broken
and Jerry must provide a
“seed” *s* such that $E = H(s)$

Is first assumption plausible?

Would the public really accept *any* curve chosen by Jerry that survives these criteria?

Example showing plausibility:
Chinese OSCCA SM2 (2010)
includes algorithms and a curve.
The curve looks random;
survives these criteria;
has no other justification.

More recent example:
French [ANSSI FRP256V1](#) (2011).
Again no justification.

Maybe public is more demanding outside China and France:
 E must not be publicly broken,
and Jerry must provide a
“seed” s such that $E = H(s)$.

Is first assumption plausible?

Would the public really accept *any* curve chosen by Jerry that survives these criteria?

Example showing plausibility:
Chinese OSCCA SM2 (2010) includes algorithms and a curve. The curve looks random; survives these criteria; has no other justification.

More recent example:
French [ANSSI FRP256V1](#) (2011).
Again no justification.

Maybe public is more demanding outside China and France:
 E must not be publicly broken, *and* Jerry must provide a “seed” s such that $E = H(s)$.

Examples: [ANSI X9.62](#) (1999) “selecting an elliptic curve verifiably at random”; [Certicom SEC 2 1.0](#) (2000) “verifiably random parameters offer some additional conservative features” — “parameters cannot be predetermined”; [NIST FIPS 186-2](#) (2000); [ANSI X9.63](#) (2001); [Certicom SEC 2 2.0](#) (2010).

assumption plausible?

the public really accept

ve chosen by Jerry

vives these criteria?

e showing plausibility:

OSCCA SM2 (2010)

algorithms and a curve.

ve looks random;

these criteria;

other justification.

cent example:

[ANSI X9.62](#) (1999).

o justification.

Maybe public is more demanding
outside China and France:

E must not be publicly broken,

and Jerry must provide a

“seed” s such that $E = H(s)$.

Examples: [ANSI X9.62](#) (1999)

“selecting an elliptic curve

verifiably at random”; [Certicom](#)

[SEC 2 1.0](#) (2000) “verifiably

random parameters offer

some additional conservative

features” — “parameters cannot

be predetermined”; [NIST FIPS](#)

[186-2](#) (2000); [ANSI X9.63](#) (2001);

[Certicom SEC 2 2.0](#) (2010).

NIST de

$y^2 = x^3$

$b^2c = -$

hash is S

plausible?

really accept

by Jerry

criteria?

plausibility:

M2 (2010)

s and a curve.

andom;

eria;

ification.

ple:

[P256V1](#) (2011).

tion.

Maybe public is more demanding
outside China and France:

E must not be publicly broken,
and Jerry must provide a
“seed” s such that $E = H(s)$.

Examples: [ANSI X9.62](#) (1999)

“selecting an elliptic curve
verifiably at random”; [Certicom](#)

[SEC 2 1.0](#) (2000) “verifiably
random parameters offer

some additional conservative
features” — “parameters cannot

be predetermined”; [NIST FIPS](#)

[186-2](#) (2000); [ANSI X9.63](#) (2001);

[Certicom SEC 2 2.0](#) (2010).

NIST defines curve

$$y^2 = x^3 - 3x + b$$

$$b^2c = -27; c \text{ is a}$$

hash is SHA-1 con

Maybe public is more demanding outside China and France:

E must not be publicly broken, *and* Jerry must provide a “seed” s such that $E = H(s)$.

Examples: [ANSI X9.62](#) (1999)

“selecting an elliptic curve verifiably at random”; [Certicom](#)

[SEC 2 1.0](#) (2000) “verifiably random parameters offer

some additional conservative features” — “parameters cannot

be predetermined”; [NIST FIPS](#)

[186-2](#) (2000); [ANSI X9.63](#) (2001);

[Certicom SEC 2 2.0](#) (2010).

NIST defines curve E as

$y^2 = x^3 - 3x + b$ where

$b^2c = -27$; c is a hash of s

hash is SHA-1 concatenation

Maybe public is more demanding outside China and France:

E must not be publicly broken, *and* Jerry must provide a “seed” s such that $E = H(s)$.

Examples: [ANSI X9.62](#) (1999)

“selecting an elliptic curve verifiably at random”; [Certicom](#)

[SEC 2 1.0](#) (2000) “verifiably random parameters offer

some additional conservative features” — “parameters cannot be predetermined”; [NIST FIPS](#)

[186-2](#) (2000); [ANSI X9.63](#) (2001);

[Certicom SEC 2 2.0](#) (2010).

NIST defines curve E as

$y^2 = x^3 - 3x + b$ where

$b^2c = -27$; c is a hash of s ;

hash is SHA-1 concatenation.

Maybe public is more demanding outside China and France:

E must not be publicly broken, *and* Jerry must provide a “seed” s such that $E = H(s)$.

Examples: [ANSI X9.62](#) (1999)

“selecting an elliptic curve verifiably at random”; [Certicom](#)

[SEC 2 1.0](#) (2000) “verifiably random parameters offer

some additional conservative features” — “parameters cannot be predetermined”; [NIST FIPS](#)

[186-2](#) (2000); [ANSI X9.63](#) (2001);

[Certicom SEC 2 2.0](#) (2010).

NIST defines curve E as

$y^2 = x^3 - 3x + b$ where

$b^2c = -27$; c is a hash of s ;

hash is SHA-1 concatenation.

1999 Scott: “Consider now the possibility that one in a million of all curves have an exploitable structure that ‘they’ know about, but we don’t. Then ‘they’ **simply generate a million random seeds** until they find one that generates one of ‘their’ curves. Then they get us to use them.”

public is more demanding
China and France:
not be publicly broken,
y must provide a
s such that $E = H(s)$.

es: [ANSI X9.62](#) (1999)
g an elliptic curve
y at random”; [Certicom](#)
.0 (2000) “verifiably
parameters offer
ditional conservative
’ — “parameters cannot
etermined”; [NIST FIPS](#)
2000); [ANSI X9.63](#) (2001);
n [SEC 2 2.0](#) (2010).

NIST defines curve E as
 $y^2 = x^3 - 3x + b$ where
 $b^2c = -27$; c is a hash of s ;
hash is SHA-1 concatenation.

1999 Scott: “Consider now the
possibility that one in a million
of all curves have an exploitable
structure that ‘they’ know about,
but we don’t. Then ‘they’ **simply**
generate a million random seeds
until they find one that generates
one of ‘their’ curves. Then they
get us to use them.”

Optimize
cluster o
 $H = Ke$
“secure-
BADA55E0
B47FCEB9
mod NIS

more demanding
France:
publicly broken,
provide a
t $E = H(s)$.
X9.62 (1999)
tic curve
m”; Certicom
“verifiably
s offer
onservative
eters cannot
; NIST FIPS
SI X9.63 (2001);
.0 (2010).

NIST defines curve E as
 $y^2 = x^3 - 3x + b$ where
 $b^2c = -27$; c is a hash of s ;
hash is SHA-1 concatenation.

1999 Scott: “Consider now the
possibility that one in a million
of all curves have an exploitable
structure that ‘they’ know about,
but we don’t. Then ‘they’ simply
generate a million random seeds
until they find one that generates
one of ‘their’ curves. Then they
get us to use them.”

Optimized this con
cluster of 41 GTX
 $H = \text{Keccak}$. In 7
“secure+twist-secr
BADA55ECD8BBEAD3
B47FCEB9BE7E0E70
mod NIST P-256.

NIST defines curve E as
 $y^2 = x^3 - 3x + b$ where
 $b^2c = -27$; c is a hash of s ;
hash is SHA-1 concatenation.

1999 Scott: "Consider now the possibility that one in a million of all curves have an exploitable structure that 'they' know about, but we don't. Then 'they' **simply generate a million random seeds** until they find one that generates one of 'their' curves. Then they get us to use them."

Optimized this computation
cluster of 41 GTX780 GPUs
 $H = \text{Keccak}$. In 7 hours for
"secure+twist-secure" $b = 0$
BADA55ECD8BBEAD3ADD6C534F
B47FCEB9BE7E0E702A8D1DD56
mod NIST P-256.

NIST defines curve E as $y^2 = x^3 - 3x + b$ where $b^2c = -27$; c is a hash of s ; hash is SHA-1 concatenation.

1999 Scott: “Consider now the possibility that one in a million of all curves have an exploitable structure that ‘they’ know about, but we don’t. Then ‘they’ **simply generate a million random seeds** until they find one that generates one of ‘their’ curves. Then they get us to use them.”

Optimized this computation on cluster of 41 GTX780 GPUs using $H = \text{Keccak}$. In 7 hours found “secure+twist-secure” $b = 0x$

BADA55ECD8BBEAD3ADD6C534F92197DE
B47FCEB9BE7E0E702A8D1DD56B5D0B0C
mod NIST P-256.

NIST defines curve E as $y^2 = x^3 - 3x + b$ where $b^2c = -27$; c is a hash of s ; hash is SHA-1 concatenation.

1999 Scott: “Consider now the possibility that one in a million of all curves have an exploitable structure that ‘they’ know about, but we don’t. Then ‘they’ **simply generate a million random seeds** until they find one that generates one of ‘their’ curves. Then they get us to use them.”

Optimized this computation on cluster of 41 GTX780 GPUs using $H = \text{Keccak}$. In 7 hours found “secure+twist-secure” $b = 0x$

BADA55ECD8BBEAD3ADD6C534F92197DE
B47FCEB9BE7E0E702A8D1DD56B5D0B0C
mod NIST P-256.

Similarly found $b = 0x$

BADA55ECFD9CA54C0738B8A6FB8CF4CC
F84E916D83D6DA1B78B622351E11AB4E
mod NIST P-224; and $b = 0x$

BADA55EC3BE2AD1F9EEEA5881ECF95BB
F3AC392526F01D4CD13E684C63A17CC4
D5F271642AD83899113817A61006413D
mod NIST P-384.

defines curve E as
 $y^2 = x^3 + b$ where
 $b = H(s) \oplus 27$; c is a hash of s ;
SHA-1 concatenation.

Scott: “Consider now the
probability that one in a million
curves have an exploitable
weakness that ‘they’ know about,
but ‘they’ don’t. Then ‘they’ simply
generate a million random seeds
and they find one that generates
a ‘weak’ curve. Then they
use them.”

Optimized this computation on
cluster of 41 GTX780 GPUs using
 $H = \text{Keccak}$. In 7 hours found
“secure+twist-secure” $b = 0x$

BADA55ECD8BBEAD3ADD6C534F92197DE
B47FCEB9BE7E0E702A8D1DD56B5D0B0C
mod NIST P-256.

Similarly found $b = 0x$

BADA55ECFD9CA54C0738B8A6FB8CF4CC
F84E916D83D6DA1B78B622351E11AB4E
mod NIST P-224; and $b = 0x$

BADA55EC3BE2AD1F9EEEA5881ECF95BB
F3AC392526F01D4CD13E684C63A17CC4
D5F271642AD83899113817A61006413D
mod NIST P-384.

Maybe i
the publ

e E as
where
hash of s ;
concatenation.

Consider now the
e in a million
an exploitable
they' know about,
en 'they' simply
random seeds
e that generates
es. Then they
n."

Optimized this computation on
cluster of 41 GTX780 GPUs using
 $H = \text{Keccak}$. In 7 hours found
"secure+twist-secure" $b = 0x$

BADA55ECD8BBEAD3ADD6C534F92197DE
B47FCEB9BE7E0E702A8D1DD56B5D0B0C
mod NIST P-256.

Similarly found $b = 0x$

BADA55ECFD9CA54C0738B8A6FB8CF4CC
F84E916D83D6DA1B78B622351E11AB4E
mod NIST P-224; and $b = 0x$

BADA55EC3BE2AD1F9EEEA5881ECF95BB
F3AC392526F01D4CD13E684C63A17CC4
D5F271642AD83899113817A61006413D
mod NIST P-384.

Maybe in some co
the public is more

Optimized this computation on
cluster of 41 GTX780 GPUs using
 $H = \text{Keccak}$. In 7 hours found

“secure+twist-secure” $b = 0x$

BADA55ECD8BBEAD3ADD6C534F92197DE
B47FCEB9BE7E0E702A8D1DD56B5D0B0C
mod NIST P-256.

Similarly found $b = 0x$

BADA55ECFD9CA54C0738B8A6FB8CF4CC
F84E916D83D6DA1B78B622351E11AB4E
mod NIST P-224; and $b = 0x$

BADA55EC3BE2AD1F9EEEA5881ECF95BB
F3AC392526F01D4CD13E684C63A17CC4
D5F271642AD83899113817A61006413D
mod NIST P-384.

Maybe in some countries
the public is more demanding

Optimized this computation on
cluster of 41 GTX780 GPUs using

$H = \text{Keccak}$. In 7 hours found

“secure+twist-secure” $b = 0x$

BADA55ECD8BBEAD3ADD6C534F92197DE

B47FCEB9BE7E0E702A8D1DD56B5D0B0C

mod NIST P-256.

Similarly found $b = 0x$

BADA55ECFD9CA54C0738B8A6FB8CF4CC

F84E916D83D6DA1B78B622351E11AB4E

mod NIST P-224; and $b = 0x$

BADA55EC3BE2AD1F9EEEA5881ECF95BB

F3AC392526F01D4CD13E684C63A17CC4

D5F271642AD83899113817A61006413D

mod NIST P-384.

Maybe in some countries
the public is more demanding.

Optimized this computation on cluster of 41 GTX780 GPUs using

$H = \text{Keccak}$. In 7 hours found

“secure+twist-secure” $b = 0x$

BADA55ECD8BBEAD3ADD6C534F92197DE
B47FCEB9BE7E0E702A8D1DD56B5D0B0C
mod NIST P-256.

Similarly found $b = 0x$

BADA55ECFD9CA54C0738B8A6FB8CF4CC
F84E916D83D6DA1B78B622351E11AB4E
mod NIST P-224; and $b = 0x$

BADA55EC3BE2AD1F9EEEA5881ECF95BB
F3AC392526F01D4CD13E684C63A17CC4
D5F271642AD83899113817A61006413D
mod NIST P-384.

Maybe in some countries the public is more demanding.

Brainpool standard (2005):

“The choice of the seeds from which the [NIST] curve parameters have been derived is not motivated leaving an essential part of the security analysis open.

... **Verifiably pseudo-random.**

The [Brainpool] curves shall be generated in a pseudo-random manner using seeds that are generated in a systematic and comprehensive way.”

ed this computation on
of 41 GTX780 GPUs using
eccak. In 7 hours found
+twist-secure" $b = 0x$
CD8BBEAD3ADD6C534F92197DE
9BE7E0E702A8D1DD56B5D0B0C
ST P-256.
y found $b = 0x$
CFD9CA54C0738B8A6FB8CF4CC
D83D6DA1B78B622351E11AB4E
ST P-224; and $b = 0x$
C3BE2AD1F9EEEA5881ECF95BB
526F01D4CD13E684C63A17CC4
42AD83899113817A61006413D
ST P-384.

Maybe in some countries
the public is more demanding.
Brainpool standard (2005):
"The choice of the seeds
from which the [NIST] curve
parameters have been derived is
not motivated leaving an essential
part of the security analysis open.
... **Verifiably pseudo-random.**
The [Brainpool] curves shall be
generated in a pseudo-random
manner using seeds that are
generated in a systematic and
comprehensive way."

```
import hashlib
def hash(seed): h
seedbytes = 20

p = 0xD7C134AA264
k = GF(p); R.<x>

def secure(A,B):
    if k(B).is_squa
    n = EllipticCur
    return (n < p a
        and Integers(

def int2str(seed,
    return ''.join(

def str2int(seed)
    return Integer(

def update(seed):
    return int2str(

def fullhash(seed
    return str2int(

def real2str(seed
    return int2str(

nums = real2str(e
S = nums[2*seedby
while True:
    A = fullhash(S)
    if not (k(A)*x^
    S = update(S)
    B = fullhash(S)
    if not secure(A
    print 'p',hex(p
    print 'A',hex(A
    print 'B',hex(B
    break
```

computation on
780 GPUs using
7 hours found
"curve" $b = 0x$
ADD6C534F92197DE
2A8D1DD56B5D0B0C

 $= 0x$
0738B8A6FB8CF4CC
78B622351E11AB4E
and $b = 0x$
9EEEA5881ECF95BB
D13E684C63A17CC4
113817A61006413D

Maybe in some countries
the public is more demanding.

Brainpool standard (2005):

"The choice of the seeds
from which the [NIST] curve
parameters have been derived is
not motivated leaving an essential
part of the security analysis open.

... **Verifiably pseudo-random.**

The [Brainpool] curves shall be
generated in a pseudo-random
manner using seeds that are
generated in a systematic and
comprehensive way."

```
import hashlib
def hash(seed): h = hashlib.sha1(); h.
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787
k = GF(p); R.<x> = k[]

def secure(A,B):
    if k(B).is_square(): return False
    n = EllipticCurve([k(A),k(B)]).cardi
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_

def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%25

def str2int(seed):
    return Integer(seed.encode('hex'),16

def update(seed):
    return int2str(str2int(seed) + 1,len

def fullhash(seed):
    return str2int(hash(seed) + hash(upd

def real2str(seed,bytes):
    return int2str(Integer(floor(RealFie

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x4+3).roots(): S = upd
    S = update(S)
    B = fullhash(S)
    if not secure(A,B): S = update(S); c
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break
```

on
using
und
x
92197DE
B5D0B0C
B8CF4CC
E11AB4E
x
ECF95BB
3A17CC4
006413D

Maybe in some countries
the public is more demanding.

Brainpool standard (2005):

“The choice of the seeds
from which the [NIST] curve
parameters have been derived is
not motivated leaving an essential
part of the security analysis open.

... **Verifiably pseudo-random.**

The [Brainpool] curves shall be
generated in a pseudo-random
manner using seeds that are
generated in a systematic and
comprehensive way.”

```
import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC80
k = GF(p); R.<x> = k[]

def secure(A,B):
    if k(B).is_square(): return False
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%256) for i in reversed(

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*2

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x4+3).roots(): S = update(S); continue
    S = update(S)
    B = fullhash(S)
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break
```

Maybe in some countries
the public is more demanding.

Brainpool standard (2005):

“The choice of the seeds
from which the [NIST] curve
parameters have been derived is
not motivated leaving an essential
part of the security analysis open.

... **Verifiably pseudo-random.**

The [Brainpool] curves shall be
generated in a pseudo-random
manner using seeds that are
generated in a systematic and
comprehensive way.”

```
import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    if k(B).is_square(): return False
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x4+3).roots(): S = update(S); continue
    S = update(S)
    B = fullhash(S)
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break
```

in some countries
is more demanding.
standard (2005):
choice of the seeds
which the [NIST] curve
parameters have been derived is
motivated leaving an essential
the security analysis open.
reliably pseudo-random.
[rainpool] curves shall be
derived in a pseudo-random
using seeds that are
derived in a systematic and
comprehensive way.”

```
import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    if k(B).is_square(): return False
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2^223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256^bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x^4+3).roots(): S = update(S); continue
    S = update(S)
    B = fullhash(S)
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break
```

We care
the curve
from the
Previous
Output

```
p D7C134AA2643
A 2B98B906DC24
B 68AEC4BFE84C
```

ountries
demanding.
d (2005):
e seeds
[IST] curve
een derived is
ving an essential
y analysis open.
pseudo-random.
urves shall be
pseudo-random
ls that are
tematic and
y.”

```
import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    if k(B).is_square(): return False
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2^223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256^bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x^4+3).roots(): S = update(S); continue
    S = update(S)
    B = fullhash(S)
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break
```

We carefully imple
the curve-generati
from the Brainpoo
Previous slide: 224
Output of this pro

```
p D7C134AA264366862A18302575D1D7
A 2B98B906DC245F2916C03A2F953EAS
B 68AEC4BFE84C659EBB8B81DC39355A
```

```

import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    if k(B).is_square(): return False
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x4+3).roots(): S = update(S); continue
    S = update(S)
    B = fullhash(S)
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break

```

We carefully implemented
the curve-generation procedure
from the Brainpool standard
Previous slide: 224-bit procedure

Output of this procedure:

```

p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 2B98B906DC245F2916C03A2F953EA9AE565C3253E8AEC4BFE84C659EBB8B81DC39355A2EBFA3870D98976FA
B 68AEC4BFE84C659EBB8B81DC39355A2EBFA3870D98976FA

```

```

import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    if k(B).is_square(): return False
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2^223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256^bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x^4+3).roots(): S = update(S); continue
    S = update(S)
    B = fullhash(S)
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break

```

We carefully implemented
the curve-generation procedure
from the Brainpool standard.
Previous slide: 224-bit procedure.

Output of this procedure:

```

p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 2B98B906DC245F2916C03A2F953EA9AE565C3253E8AEC4BFE84C659E
B 68AEC4BFE84C659EBB8B81DC39355A2EBFA3870D98976FA2F17D2D8D

```



```

import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    if k(B).is_square(): return False
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2^223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256^bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x^4+3).roots(): S = update(S); continue
    S = update(S)
    B = fullhash(S)
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break

```

We carefully implemented
the curve-generation procedure
from the Brainpool standard.
Previous slide: 224-bit procedure.

Output of this procedure:

```

p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 2B98B906DC245F2916C03A2F953EA9AE565C3253E8AEC4BFE84C659E
B 68AEC4BFE84C659EBB8B81DC39355A2EBFA3870D98976FA2F17D2D8D

```

The standard 224-bit Brainpool
curve **is not the same curve:**

```

p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 68A5E62CA9CE6C1C299803A6C1530B514E182AD8B0042A59CAD29F43
B 2580F63CCFE44138870713B1A92369E33E2135D266DBB372386C400B

```

```

import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    if k(B).is_square(): return False
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2^223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256^bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x^4+3).roots(): S = update(S); continue
    S = update(S)
    B = fullhash(S)
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break

```

We carefully implemented
the curve-generation procedure
from the Brainpool standard.
Previous slide: 224-bit procedure.

Output of this procedure:

```

p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 2B98B906DC245F2916C03A2F953EA9AE565C3253E8AEC4BFE84C659E
B 68AEC4BFE84C659EBB8B81DC39355A2EBFA3870D98976FA2F17D2D8D

```

The standard 224-bit Brainpool
curve **is not the same curve:**

```

p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 68A5E62CA9CE6C1C299803A6C1530B514E182AD8B0042A59CAD29F43
B 2580F63CCFE44138870713B1A92369E33E2135D266DBB372386C400B

```

Next slide: a procedure
that **does** generate
the standard Brainpool curve.

```

= hashlib.sha1(); h.update(seed); return h.digest()

366862A18302575D1D787B09F075797DA89F57EC8C0FF
= k[]

re(): return False
ve([k(A),k(B)]).cardinality()
nd n.is_prime()
n(p).multiplicative_order() * 100 >= n-1)

bytes):
[chr((seed//256^i)%256) for i in reversed(range(bytes))]]

seed.encode('hex'),16)

str2int(seed) + 1,len(seed))

):
hash(seed) + hash(update(seed))) % 2^223

,bytes):
Integer(floor(RealField(8*bytes+8)(seed)*256^bytes)),bytes)

xp(1)/16,7*seedbytes)
tes:3*seedbytes]

4+3).roots(): S = update(S); continue

,B): S = update(S); continue
).upper()
).upper()
).upper()

```

We carefully implemented
the curve-generation procedure
from the Brainpool standard.
Previous slide: 224-bit procedure.

Output of this procedure:

```

p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 2B98B906DC245F2916C03A2F953EA9AE565C3253E8AEC4BFE84C659E
B 68AEC4BFE84C659EBB8B81DC39355A2EBFA3870D98976FA2F17D2D8D

```

The standard 224-bit Brainpool
curve **is not the same curve:**

```

p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 68A5E62CA9CE6C1C299803A6C1530B514E182AD8B0042A59CAD29F43
B 2580F63CCFE44138870713B1A92369E33E2135D266DBB372386C400B

```

Next slide: a procedure
that **does** generate
the standard Brainpool curve.

```

import hashlib
def hash(seed): h
seedbytes = 20

p = 0xD7C134AA264
k = GF(p); R.<x>

def secure(A,B):
    n = EllipticCur
    return (n < p and
            and Integers(

def int2str(seed,
    return ''.join(

def str2int(seed)
    return Integer(

def update(seed):
    return int2str(

def fullhash(seed)
    return str2int(

def real2str(seed)
    return int2str(

nums = real2str(e
S = nums[2*seedby
while True:
    A = fullhash(S)
    if not (k(A)*x^
        while True:
            S = update(S)
            B = fullhash(
                if not k(B).i
    if not secure(A
    print 'p',hex(p
    print 'A',hex(A
    print 'B',hex(B
    break

```

```
update(seed); return h.digest()
```

```
B09F075797DA89F57EC8C0FF
```

```
nality()
```

```
order() * 100 >= n-1)
```

```
6) for i in reversed(range(bytes))])
```

```
)
```

```
(seed))
```

```
ate(seed))) % 2223
```

```
ld(8*bytes+8)(seed)*256bytes),bytes)
```

```
ate(S); continue
```

```
ontinue
```

We carefully implemented
the curve-generation procedure
from the Brainpool standard.
Previous slide: 224-bit procedure.

Output of this procedure:

```
p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 2B98B906DC245F2916C03A2F953EA9AE565C3253E8AEC4BFE84C659E
B 68AEC4BFE84C659EBB8B81DC39355A2EBFA3870D98976FA2F17D2D8D
```

The standard 224-bit Brainpool
curve **is not the same curve:**

```
p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 68A5E62CA9CE6C1C299803A6C1530B514E182AD8B0042A59CAD29F43
B 2580F63CCFE44138870713B1A92369E33E2135D266DBB372386C400B
```

Next slide: a procedure
that **does** generate
the standard Brainpool curve.

```
import hashlib
def hash(seed): h = hashlib.sha1(); h.
seedbytes = 20
```

```
p = 0xD7C134AA264366862A18302575D1D787
k = GF(p); R.<x> = k[]
```

```
def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardi
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_
```

```
def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%25
```

```
def str2int(seed):
    return Integer(seed.encode('hex'),16
```

```
def update(seed):
    return int2str(str2int(seed) + 1,len
```

```
def fullhash(seed):
    return str2int(hash(seed) + hash(upd
```

```
def real2str(seed,bytes):
    return int2str(Integer(floor(RealFie
```

```
nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
```

```
while True:
    A = fullhash(S)
    if not (k(A)*x4+3).roots(): S = upd
    while True:
        S = update(S)
        B = fullhash(S)
        if not k(B).is_square(): break
    if not secure(A,B): S = update(S); c
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break
```

We carefully implemented
the curve-generation procedure
from the Brainpool standard.
Previous slide: 224-bit procedure.

Output of this procedure:

```
p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 2B98B906DC245F2916C03A2F953EA9AE565C3253E8AEC4BFE84C659E
B 68AEC4BFE84C659EBB8B81DC39355A2EBFA3870D98976FA2F17D2D8D
```

The standard 224-bit Brainpool
curve is not the same curve:

```
p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 68A5E62CA9CE6C1C299803A6C1530B514E182AD8B0042A59CAD29F43
B 2580F63CCFE44138870713B1A92369E33E2135D266DBB372386C400B
```

Next slide: a procedure
that **does** generate
the standard Brainpool curve.

```
import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x4+3).roots(): S = update(S); continue
    while True:
        S = update(S)
        B = fullhash(S)
        if not k(B).is_square(): break
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break
```

We carefully implemented
the curve-generation procedure
from the Brainpool standard.
Previous slide: 224-bit procedure.

Output of this procedure:

```
p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 2B98B906DC245F2916C03A2F953EA9AE565C3253E8AEC4BFE84C659E
B 68AEC4BFE84C659EBB8B81DC39355A2EBFA3870D98976FA2F17D2D8D
```

The standard 224-bit Brainpool
curve **is not the same curve:**

```
p D7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
A 68A5E62CA9CE6C1C299803A6C1530B514E182AD8B0042A59CAD29F43
B 2580F63CCFE44138870713B1A92369E33E2135D266DBB372386C400B
```

Next slide: a procedure
that **does** generate
the standard Brainpool curve.

```
import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2^223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256^bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x^4+3).roots(): S = update(S); continue
    while True:
        S = update(S)
        B = fullhash(S)
        if not k(B).is_square(): break
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break
```

fully implemented
e-generation procedure
e Brainpool standard.
s slide: 224-bit procedure.

of this procedure:

66862A18302575D1D787B09F075797DA89F57EC8C0FF
5F2916C03A2F953EA9AE565C3253E8AEC4BFE84C659E
659EBB8B81DC39355A2EBFA3870D98976FA2F17D2D8D

standard 224-bit Brainpool
not the same curve:

66862A18302575D1D787B09F075797DA89F57EC8C0FF
6C1C299803A6C1530B514E182AD8B0042A59CAD29F43
4138870713B1A92369E33E2135D266DBB372386C400B

de: a procedure
es generate
dard Brainpool curve.

```
import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x4+3).roots(): S = update(S); continue
    while True:
        S = update(S)
        B = fullhash(S)
        if not k(B).is_square(): break
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break
```

Did Brain
publicati
Did they
Brainpoo
advertise
“compre
transpar
say the s
Can quic
to take t

plemented
on procedure
ol standard.
4-bit procedure.

cedure:

87B09F075797DA89F57EC8C0FF
AE565C3253E8AEC4BFE84C659E
2EBFA3870D98976FA2F17D2D8D

bit Brainpool
ame curve:

87B09F075797DA89F57EC8C0FF
3514E182AD8B0042A59CAD29F43
E33E2135D266DBB372386C400B

edure

e
npool curve.

```
import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2^223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256^bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x^4+3).roots(): S = update(S); continue
    while True:
        S = update(S)
        B = fullhash(S)
        if not k(B).is_square(): break
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break
```

Did Brainpool check
publication? After
Did they know before
Brainpool procedure
advertised as “system
“comprehensive”,
transparent”, etc.
say the same for *b*
Can quietly manipulate
to take the weaker


```

import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x4+3).roots(): S = update(S); continue
    while True:
        S = update(S)
        B = fullhash(S)
        if not k(B).is_square(): break
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break

```

Did Brainpool check before publication? After publication Did they know before 2015?

Brainpool procedure is advertised as “systematic”, “comprehensive”, “complete transparent”, etc. Surely we say the same for *both* procedures.

Can quietly manipulate choices to take the weaker procedure?

```

import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2^223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256^bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x^4+3).roots(): S = update(S); continue
    while True:
        S = update(S)
        B = fullhash(S)
        if not k(B).is_square(): break
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break

```

Did Brainpool check before
publication? After publication?

Did they know before 2015?

Brainpool procedure is
advertised as “systematic”,
“comprehensive”, “completely
transparent”, etc. Surely we can
say the same for *both* procedures.

Can quietly manipulate choice
to take the weaker procedure.

```

import hashlib
def hash(seed): h = hashlib.sha1(); h.update(seed); return h.digest()
seedbytes = 20

p = 0xD7C134AA264366862A18302575D1D787B09F075797DA89F57EC8C0FF
k = GF(p); R.<x> = k[]

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n < p and n.is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def update(seed):
    return int2str(str2int(seed) + 1,len(seed))

def fullhash(seed):
    return str2int(hash(seed) + hash(update(seed))) % 2^223

def real2str(seed,bytes):
    return int2str(Integer(floor(RealField(8*bytes+8)(seed)*256^bytes)),bytes)

nums = real2str(exp(1)/16,7*seedbytes)
S = nums[2*seedbytes:3*seedbytes]
while True:
    A = fullhash(S)
    if not (k(A)*x^4+3).roots(): S = update(S); continue
    while True:
        S = update(S)
        B = fullhash(S)
        if not k(B).is_square(): break
    if not secure(A,B): S = update(S); continue
    print 'p',hex(p).upper()
    print 'A',hex(A).upper()
    print 'B',hex(B).upper()
    break

```

Did Brainpool check before publication? After publication?

Did they know before 2015?

Brainpool procedure is advertised as “systematic”, “comprehensive”, “completely transparent”, etc. Surely we can say the same for *both* procedures.

Can quietly manipulate choice to take the weaker procedure.

Interesting Brainpool quote: “It is envisioned to provide additional curves on a regular basis.”

```

= hashlib.sha1(); h.update(seed); return h.digest()

366862A18302575D1D787B09F075797DA89F57EC8C0FF
= k[]

ve([k(A),k(B)]).cardinality()
ndn.is_prime()
n(p).multiplicative_order() * 100 >= n-1

bytes):
[chr((seed//256^i)%256) for i in reversed(range(bytes))]]

:
seed.encode('hex'),16)

str2int(seed) + 1,len(seed))

):
hash(seed) + hash(update(seed))) % 2^223

,bytes):
Integer(floor(RealField(8*bytes+8)(seed)*256^bytes)),bytes)

xp(1)/16,7*seedbytes)
tes:3*seedbytes]

4+3).roots(): S = update(S); continue

S)
s_square(): break
,B): S = update(S); continue
).upper()
).upper()
).upper()

```

Did Brainpool check before publication? After publication?

Did they know before 2015?

Brainpool procedure is advertised as “systematic”, “comprehensive”, “completely transparent”, etc. Surely we can say the same for *both* procedures.

Can quietly manipulate choice to take the weaker procedure.

Interesting Brainpool quote: “It is envisioned to provide additional curves on a regular basis.”

We made using sta

To avoid

complica

hash out

from SH

maximum

Also upg

maximum

Brainpoo

and arct

uses sin(

We also

pattern

```
update(seed); return h.digest()
```

```
B09F075797DA89F57EC8C0FF
```

```
inality()
```

```
order() * 100 >= n-1)
```

```
6) for i in reversed(range(bytes))])
```

```
(seed))
```

```
ate(seed))) % 2^223
```

```
ld(8*bytes+8)(seed)*256^bytes),bytes)
```

```
ate(S); continue
```

```
ontinue
```

Did Brainpool check before publication? After publication?

Did they know before 2015?

Brainpool procedure is advertised as “systematic”, “comprehensive”, “completely transparent”, etc. Surely we can say the same for *both* procedures.

Can quietly manipulate choice to take the weaker procedure.

Interesting Brainpool quote: “It is envisioned to provide additional curves on a regular basis.”

We made a new 2 using standard NIS

To avoid Brainpool complications of c hash outputs: We from SHA-1 to sta maximum-security Also upgraded to maximum twist se

Brainpool uses exp and $\arctan(1) = \pi/4$ uses $\sin(1)$, so we We also used muc pattern of searching

Did Brainpool check before publication? After publication? Did they know before 2015?

Brainpool procedure is advertised as “systematic”, “comprehensive”, “completely transparent”, etc. Surely we can say the same for *both* procedures.

Can quietly manipulate choice to take the weaker procedure.

Interesting Brainpool quote: “It is envisioned to provide additional curves on a regular basis.”

We made a new 224-bit curve using standard NIST P-224

To avoid Brainpool’s complications of concatenating hash outputs: We upgraded from SHA-1 to state-of-the-art maximum-security SHA3-512. Also upgraded to requiring maximum twist security.

Brainpool uses $\exp(1) = e$ and $\arctan(1) = \pi/4$, and M uses $\sin(1)$, so we used $\cos(1)$. We also used much simpler pattern of searching for seeds.

Did Brainpool check before publication? After publication?

Did they know before 2015?

Brainpool procedure is advertised as “systematic”, “comprehensive”, “completely transparent”, etc. Surely we can say the same for *both* procedures.

Can quietly manipulate choice to take the weaker procedure.

Interesting Brainpool quote: “It is envisioned to provide additional curves on a regular basis.”

We made a new 224-bit curve using standard NIST P-224 prime.

To avoid Brainpool’s complications of concatenating hash outputs: We upgraded from SHA-1 to state-of-the-art maximum-security SHA3-512.

Also upgraded to requiring maximum twist security.

Brainpool uses $\exp(1) = e$ and $\arctan(1) = \pi/4$, and MD5 uses $\sin(1)$, so we used $\cos(1)$.

We also used much simpler pattern of searching for seeds.

inpool check before
ion? After publication?
y know before 2015?
ol procedure is
ed as “systematic”,
prehensive”, “completely
ent”, etc. Surely we can
same for *both* procedures.
etly manipulate choice
the weaker procedure.
ng Brainpool quote: “It
oned to provide additional
n a regular basis.”

We made a new 224-bit curve
using standard NIST P-224 prime.

To avoid Brainpool’s
complications of concatenating
hash outputs: We upgraded
from SHA-1 to state-of-the-art
maximum-security SHA3-512.

Also upgraded to requiring
maximum twist security.

Brainpool uses $\exp(1) = e$
and $\arctan(1) = \pi/4$, and MD5
uses $\sin(1)$, so we used $\cos(1)$.

We also used much simpler
pattern of searching for seeds.

```
import simplesha3
hash = simplesha3

p = 2^224 - 2^96
k = GF(p)
seedbytes = 20

def secure(A,B):
    n = EllipticCur
    return (n.is_pr
        and Integers(
        and Integers(

def int2str(seed,
    return ''.join(

def str2int(seed)
    return Integer(

def complement(se
    return ''.join(

def real2str(seed
    return int2str(

sizeofint = 4
nums = real2str(c
for counter in xr
    S = int2str(cou
    T = complement(
    A = str2int(has
    B = str2int(has
    if secure(A,B):
        print 'p',hex
        print 'A',hex
        print 'B',hex
        break
```


ck before
r publication?
fore 2015?
re is
tematic",
"completely
Surely we can
both procedures.
ulate choice
r procedure.
ool quote: "It
ovide additional
r basis."

We made a new 224-bit curve
using standard NIST P-224 prime.

To avoid Brainpool's
complications of concatenating
hash outputs: We upgraded
from SHA-1 to state-of-the-art
maximum-security SHA3-512.

Also upgraded to requiring
maximum twist security.

Brainpool uses $\exp(1) = e$
and $\arctan(1) = \pi/4$, and MD5
uses $\sin(1)$, so we used $\cos(1)$.
We also used much simpler
pattern of searching for seeds.

```
import simplesha3
hash = simplesha3.sha3512

p = 2^224 - 2^96 + 1
k = GF(p)
seedbytes = 20

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardi
    return (n.is_prime() and (2*p+2-n).i
            and Integers(n)(p).multiplicative_
            and Integers(2*p+2-n)(p).multiplic

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in range(bytes)])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def complement(seed):
    return ''.join([chr(255-ord(s)) for s in seed])

def real2str(seed,bytes):
    return int2str(Integer(RealField(8*b

sizeofint = 4
nums = real2str(cos(1),seedbytes - siz
for counter in xrange(0,256^sizeofint)
    S = int2str(counter,sizeofint) + num
    T = complement(S)
    A = str2int(hash(S))
    B = str2int(hash(T))
    if secure(A,B):
        print 'p',hex(p).upper()
        print 'A',hex(A).upper()
        print 'B',hex(B).upper()
        break
```

We made a new 224-bit curve using standard NIST P-224 prime.

To avoid Brainpool's complications of concatenating hash outputs: We upgraded from SHA-1 to state-of-the-art maximum-security SHA3-512.

Also upgraded to requiring maximum twist security.

Brainpool uses $\exp(1) = e$ and $\arctan(1) = \pi/4$, and MD5 uses $\sin(1)$, so we used $\cos(1)$. We also used much simpler pattern of searching for seeds.

```
import simplesha3
hash = simplesha3.sha3512

p = 2^224 - 2^96 + 1
k = GF(p)
seedbytes = 20

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n.is_prime() and (2*p+2-n).is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1
            and Integers(2*p+2-n)(p).multiplicative_order() * 100 >

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def complement(seed):
    return ''.join([chr(255-ord(s)) for s in seed])

def real2str(seed,bytes):
    return int2str(Integer(RealField(8*bytes)(seed)*256^bytes

sizeofint = 4
nums = real2str(cos(1),seedbytes - sizeofint)
for counter in xrange(0,256^sizeofint):
    S = int2str(counter,sizeofint) + nums
    T = complement(S)
    A = str2int(hash(S))
    B = str2int(hash(T))
    if secure(A,B):
        print 'p',hex(p).upper()
        print 'A',hex(A).upper()
        print 'B',hex(B).upper()
        break
```

We made a new 224-bit curve using standard NIST P-224 prime.

To avoid Brainpool's complications of concatenating hash outputs: We upgraded from SHA-1 to state-of-the-art maximum-security SHA3-512.

Also upgraded to requiring maximum twist security.

Brainpool uses $\exp(1) = e$ and $\arctan(1) = \pi/4$, and MD5 uses $\sin(1)$, so we used $\cos(1)$.

We also used much simpler pattern of searching for seeds.

```
import simplesha3
hash = simplesha3.sha3512

p = 2224 - 296 + 1
k = GF(p)
seedbytes = 20

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n.is_prime() and (2*p+2-n).is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1
            and Integers(2*p+2-n)(p).multiplicative_order() * 100 >= 2*p+2-n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def complement(seed):
    return ''.join([chr(255-ord(s)) for s in seed])

def real2str(seed,bytes):
    return int2str(Integer(RealField(8*bytes)(seed)*256bytes),bytes)

sizeofint = 4
nums = real2str(cos(1),seedbytes - sizeofint)
for counter in xrange(0,256sizeofint):
    S = int2str(counter,sizeofint) + nums
    T = complement(S)
    A = str2int(hash(S))
    B = str2int(hash(T))
    if secure(A,B):
        print 'p',hex(p).upper()
        print 'A',hex(A).upper()
        print 'B',hex(B).upper()
        break
```

...e a new 224-bit curve
...standard NIST P-224 prime.

...d Brainpool's

...ations of concatenating

...puts: We upgraded

...A-1 to state-of-the-art

...m-security SHA3-512.

...graded to requiring

...m twist security.

...ol uses $\exp(1) = e$

...an(1) = $\pi/4$, and MD5

... (1), so we used $\cos(1)$.

...used much simpler

...of searching for seeds.

```
import simplesha3
hash = simplesha3.sha3512

p = 2^224 - 2^96 + 1
k = GF(p)
seedbytes = 20

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n.is_prime() and (2*p+2-n).is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1
            and Integers(2*p+2-n)(p).multiplicative_order() * 100 >= 2*p+2-n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def complement(seed):
    return ''.join([chr(255-ord(s)) for s in seed])

def real2str(seed,bytes):
    return int2str(Integer(RealField(8*bytes)(seed)*256^bytes),bytes)

sizeofint = 4
nums = real2str(cos(1),seedbytes - sizeofint)
for counter in xrange(0,256^sizeofint):
    S = int2str(counter,sizeofint) + nums
    T = complement(S)
    A = str2int(hash(S))
    B = str2int(hash(T))
    if secure(A,B):
        print 'p',hex(p).upper()
        print 'A',hex(A).upper()
        print 'B',hex(B).upper()
        break
```

Output:

24-bit curve
ST P-224 prime.

ol's

concatenating

upgraded

state-of-the-art

SHA3-512.

requiring

curity.

$\cos(1) = e$

$\pi/4$, and MD5

used $\cos(1)$.

h simpler

ng for seeds.

```
import simplesha3
hash = simplesha3.sha3512

p = 2^224 - 2^96 + 1
k = GF(p)
seedbytes = 20

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n.is_prime() and (2*p+2-n).is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1
            and Integers(2*p+2-n)(p).multiplicative_order() * 100 >= 2*p+2-n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def complement(seed):
    return ''.join([chr(255-ord(s)) for s in seed])

def real2str(seed,bytes):
    return int2str(Integer(RealField(8*bytes)(seed)*256^bytes),bytes)

sizeofint = 4
nums = real2str(cos(1),seedbytes - sizeofint)
for counter in xrange(0,256^sizeofint):
    S = int2str(counter,sizeofint) + nums
    T = complement(S)
    A = str2int(hash(S))
    B = str2int(hash(T))
    if secure(A,B):
        print 'p',hex(p).upper()
        print 'A',hex(A).upper()
        print 'B',hex(B).upper()
        break
```

Output: 7144BA12CE8A

ve
prime.

ing

art

2.

MD5

1).

ls.

```
import simplesha3
hash = simplesha3.sha3512

p = 2^224 - 2^96 + 1
k = GF(p)
seedbytes = 20

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n.is_prime() and (2*p+2-n).is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1
            and Integers(2*p+2-n)(p).multiplicative_order() * 100 >= 2*p+2-n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256^i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def complement(seed):
    return ''.join([chr(255-ord(s)) for s in seed])

def real2str(seed,bytes):
    return int2str(Integer(RealField(8*bytes)(seed)*256^bytes),bytes)

sizeofint = 4
nums = real2str(cos(1),seedbytes - sizeofint)
for counter in xrange(0,256^sizeofint):
    S = int2str(counter,sizeofint) + nums
    T = complement(S)
    A = str2int(hash(S))
    B = str2int(hash(T))
    if secure(A,B):
        print 'p',hex(p).upper()
        print 'A',hex(A).upper()
        print 'B',hex(B).upper()
        break
```

Output: 7144BA12CE8A0C3BEFA053EDB

```

import simplesha3
hash = simplesha3.sha3512

p = 2224 - 296 + 1
k = GF(p)
seedbytes = 20

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n.is_prime() and (2*p+2-n).is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1
            and Integers(2*p+2-n)(p).multiplicative_order() * 100 >= 2*p+2-n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def complement(seed):
    return ''.join([chr(255-ord(s)) for s in seed])

def real2str(seed,bytes):
    return int2str(Integer(RealField(8*bytes)(seed)*256bytes),bytes)

sizeofint = 4
nums = real2str(cos(1),seedbytes - sizeofint)
for counter in xrange(0,256sizeofint):
    S = int2str(counter,sizeofint) + nums
    T = complement(S)
    A = str2int(hash(S))
    B = str2int(hash(T))
    if secure(A,B):
        print 'p',hex(p).upper()
        print 'A',hex(A).upper()
        print 'B',hex(B).upper()
        break

```

Output: 7144BA12CE8A0C3BEFA053EDBADA55...

```

import simplesha3
hash = simplesha3.sha3512

p = 2224 - 296 + 1
k = GF(p)
seedbytes = 20

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n.is_prime() and (2*p+2-n).is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1
            and Integers(2*p+2-n)(p).multiplicative_order() * 100 >= 2*p+2-n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def complement(seed):
    return ''.join([chr(255-ord(s)) for s in seed])

def real2str(seed,bytes):
    return int2str(Integer(RealField(8*bytes)(seed)*256bytes),bytes)

sizeofint = 4
nums = real2str(cos(1),seedbytes - sizeofint)
for counter in xrange(0,256sizeofint):
    S = int2str(counter,sizeofint) + nums
    T = complement(S)
    A = str2int(hash(S))
    B = str2int(hash(T))
    if secure(A,B):
        print 'p',hex(p).upper()
        print 'A',hex(A).upper()
        print 'B',hex(B).upper()
        break

```

Output: 7144BA12CE8A0C3BEFA053EDBADA55...

We actually generated >1000000 curves for this prime, each having a Brainpool-like explanation, even without complicating hashing, seed search, etc.; e.g., BADA55-VPR2-224 uses $\exp(1)$.


```

import simplesha3
hash = simplesha3.sha3512

p = 2224 - 296 + 1
k = GF(p)
seedbytes = 20

def secure(A,B):
    n = EllipticCurve([k(A),k(B)]).cardinality()
    return (n.is_prime() and (2*p+2-n).is_prime()
            and Integers(n)(p).multiplicative_order() * 100 >= n-1
            and Integers(2*p+2-n)(p).multiplicative_order() * 100 >= 2*p+2-n-1)

def int2str(seed,bytes):
    return ''.join([chr((seed//256i)%256) for i in reversed(range(bytes))])

def str2int(seed):
    return Integer(seed.encode('hex'),16)

def complement(seed):
    return ''.join([chr(255-ord(s)) for s in seed])

def real2str(seed,bytes):
    return int2str(Integer(RealField(8*bytes)(seed)*256bytes),bytes)

sizeofint = 4
nums = real2str(cos(1),seedbytes - sizeofint)
for counter in xrange(0,256sizeofint):
    S = int2str(counter,sizeofint) + nums
    T = complement(S)
    A = str2int(hash(S))
    B = str2int(hash(T))
    if secure(A,B):
        print 'p',hex(p).upper()
        print 'A',hex(A).upper()
        print 'B',hex(B).upper()
        break

```

Output: 7144BA12CE8A0C3BEFA053ED**BADA55**...

We actually generated >1000000 curves for this prime, each having a Brainpool-like explanation, even without complicating hashing, seed search, etc.; e.g., BADA55-VPR2-224 uses $\exp(1)$.

See bada55.cr.yp.to for much more: full paper; scripts; detailed Brainpool analysis; manipulating “minimal” primes and curves (Microsoft “NUMS”); manipulating security criteria.